#### Power to the points: Local certificates for clustering



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# **Data Mining Pipeline**



Learning algorithms ensure (global) quality of inference process

But what about the (local) labels assigned to data ?

Can we find LOCAL and SUCCINCT certificates that validate correctness of data labels ?

#### Local Validation in Classification



Platt scaling (P99):  $p(y = 1 | \mathbf{x}, \mathbf{w}, b) = \frac{1}{\exp(A(\langle \mathbf{w}, \mathbf{x} \rangle + b) + B)}$ 

#### Parameters are estimated using ML

# **Clustering Data**



<u>Group objects into meaningful clusters</u>

Different methods produce different answers

 k-means/medoids, HAC, spectral clustering, subspace clustering, correlation clustering, information bottleneck, ...

How do we know if an answer is good ?

# Validating Clusterings



Relative validation/ stability: compare different runs of algorithm



# Power to the points

Given a clustering of data, determine confidence scores for the label assigned to a point.

Desiderata:

- 1. Data-independent scale.
- 2. Agnostic to the method by which the clustering was made.
- 3. Works for a single clustering...
- 4. but can be used to compare different clusterings.

#### **Outlier Detection vs Local Validation**



 $\min_{S \subset P, |S| \ge (1-\epsilon)|P|} \min_{\mathcal{C}(S)} f(\mathcal{C})$ 

Outlier changes cost function but not the structure of the answer



Locally unstable points change the structure of the answer, but not the cost

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Distances are "one-dimensional" measures of influence

#### Regions can be shielded



q is equidistant from C and C' (and half the distance from C\*), and by distance estimation alone should have same chance of being assigned to either as to C\*

#### Voronoi property of clusterings



It's always better to assign a point to its nearest neighbor







Affinity of a point for a cluster is the fractional area stolen from it

$$\boldsymbol{\alpha}(p) = (\alpha_1, \dots, \alpha_k)$$
$$\sum \alpha_i = 1$$



#### Affinity of a point for a cluster is the fractional area stolen from it

- A point is "stable" if the maximum affinity is more than 0.5:
- Maximum affinity is a continuous scalar function
- This idea was first used for doing interpolation of a scalar field (natural neighbor interpolation)



#### Incorporating cluster density

If Voronoi diagram has polyhedral cells, then all relevant volumes are polyhedral cells.





$$d(p, x) = \|p - x\|^2 - w_x$$

(power diagram)



#### Generalizing to other distance spaces

Bregman divergences: Kullback-Leibler, Itakura-Saito, ...

$$B_{\phi}(x,y) = \phi(x) - \phi(y) - \langle \nabla \phi(y), x - y \rangle$$
$$d(p,x) = B_{\phi}(p,x) - w_x$$
$$d(\mathbf{p},x) = d(\mathbf{p},y) \equiv c + \langle \nabla \phi(y) - \nabla \phi(x), \mathbf{p} \rangle = 0$$

Kernel distances: graphs, strings, ...

$$d(p,x) = \|\Phi(p) - \Phi(x)\|^2 - w_x$$

# **Computing affinity vectors**

In 2D:

- Computing Voronoi diagram is O(k log k)
- Intersection of two convex polygons takes O(k) time
- k-vertex polygon can be triangulated in O(k) time
- Area of a triangle can be computed in O(1) time.

Overall: O(k log k) time per query

In 3D:

- Voronoi diagram takes O(k<sup>2</sup>) time.
- Intersection of convex polyhedra takes O(k) time
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Overall:  $O(k^2)$  time per query

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#### **Approximate Affinities**

Given  $\varepsilon > 0$  find  $\tilde{\alpha}$  such that  $|\tilde{\alpha} - \alpha| \leq \varepsilon$ 



Sampling algorithm:

- Sample s from Voronoi cell of p
- Find second closest neighbor of s
- Increment count of that neighbor
- At end, return normalized counts.

Each sample is processed in O(k) time Need to solve two problems:

- 1) How many samples to pick
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- Voronoi cell of p is a convex body
- Membership oracle is easy: "is sample nearer to p than to any other point"
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  O\*(d<sup>4</sup>) samples suffice [LV06].
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#### **Dimensionality Reduction**

Running time is polynomial in d (ambient space dimension).

Consider Euclidean distance:

$$d(x,y) = \|x - y\|^2$$

k clusters induce a k-1 dimensional space  $\mathcal{H}$ 

$$x = u + w, u \in \mathcal{H}, w \perp u$$
$$d(x, x') = \|u - u'\|^2 + \|w - w'\|^2$$

Any Voronoi cell can be written as

$$V=V'+\mathcal{H}^{\perp}$$
 ,  $V'\in\mathcal{H}$ 

Volume ratios need only be measured in  $\mathcal{H}$ 

## Algorithm Summary

Given k clusters and query p

- Project clusters to k-dimensional space
- Sample uniformly from Voronoi cell of p
- Compute frequencies of second-nearest neighbors
- Return approximate affinity scores

Overall running time:  $poly(k, 1/\epsilon)$ In practice: on the order of milliseconds/query.

# **Clustering digits**

#### Highly stable points



Highly unstable points

5 10

15

20 25

5 10 15

20 25

25

20

5

10 15

5 10

15 20 25

# **Accelerating Clustering**

"Active clustering": only pick points that inform true decision boundary

Idea: use affinity scores to identify points that might lie on boundary

- Use fast procedure to generate cluster centers (k-means++ initialization)
- Sample points with low affinity scores, as well as few points with high affinity scores.
- Cluster reduced sample.

Result: comparable accuracy of clustering with orders of magnitude speedup

Work in progress (with Kilian Weinberger): speeding up classification algorithms using affinity scores.

Points with low affinity scores act as sparse "skeleton" of data set.

#### **Model Selection**

How do we choose the right "k" for a clustering with k centers ?



#### **Model Selection**

Average stability does not increase monotonically with increasing clusters





Can affinity scores be correlated with the probabilities extracted from a clustering model ?

The (maximum) affinities define a (scalar) field over the data. Can topological methods like persistence help to identify "interesting" parts of the space ?

Can we compute points of low affinity (the data skeleton) quickly (without exploring the entire space) ?

Are there other applications where affinity scores can be used as an accelerant ?