

Limited Commitment: Mechanism Design meets Information Design

Data-Driven Decision Processes Bootcamp

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Based on joint work with Vasiliki Skreta



(Static) Mechanism Design:

- Agents have private information: T_i is the set of types of agent i and

$$\psi : \Theta \mapsto \Delta(T_1 \times \dots \times T_N)$$

describes the information player i has about θ and the types of other players.

- Payoffs only depend on A_0 .
- We are given a mapping $\pi : \Theta \mapsto \Delta(A_0)$.
- **Question:** Can we design actions for each player A_1, \dots, A_N and an outcome function $f : \times_{i=1}^N A_i \mapsto \Delta(A_0)$ such that π is the equilibrium outcome of the game defined by $\langle G, \psi, f \rangle$?

Example: Google ad auction design

- $A_0 \subseteq (\{0, 1\} \times \mathbb{R})^N$ and $(q, t) \in A_0$ if, and only if, $0 \leq \sum_{i=1}^N q_i \leq 1$.
- $\Theta = \Theta_1 \times \dots \times \Theta_N$; $T_i = \Theta_i$ denotes advertiser i 's value for the slot; $\psi(\cdot | \theta) = \delta_\theta$.
- π is the rule that assigns the good to the advertiser w/ highest θ_i .

(Static) Mechanism Design: (in more standard textbook notation)

- Agents have private information: $\Theta = \times_{i=1}^N \Theta_i$ and agent i knows θ_i . That is,

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- We need to be able to consider *all* possible games
- Mechanism design provides us with a language to do this via the **revelation principle**.

Theorem (Gibbard, 1973; Myerson, 1979; Dasgupta et al, 1979)

There exists a game that has π as an equilibrium outcome if and only if the following game implements π :

1. *Actions $M_i = \Theta_i$*
2. *When players take actions $\theta' = (\theta'_1, \dots, \theta'_N)$, the outcome is $f(\bar{\theta}) = \pi(\cdot | \theta')$.*

Furthermore, it is without loss of generality to assume that the players find it optimal to truthfully report their types.

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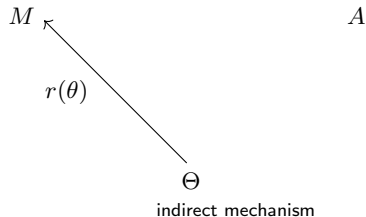
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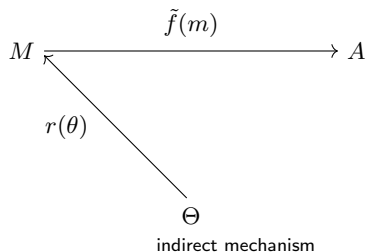


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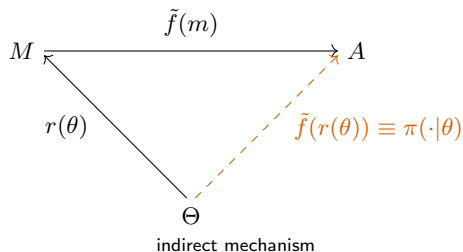


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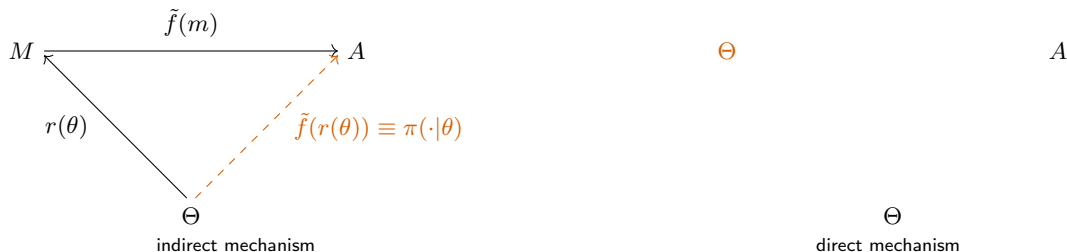


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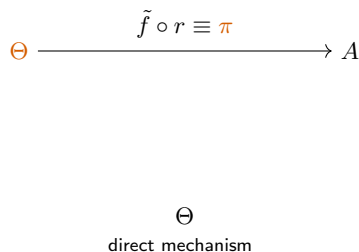
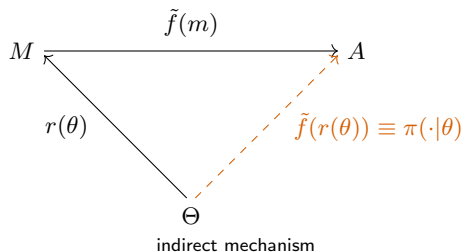


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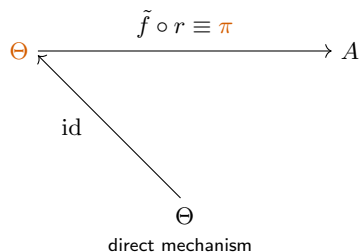
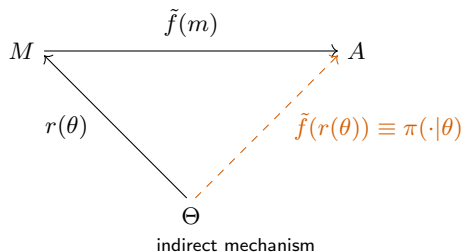


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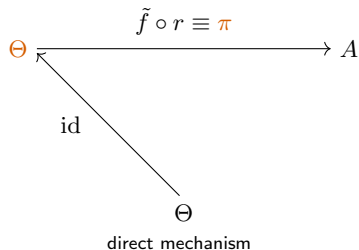
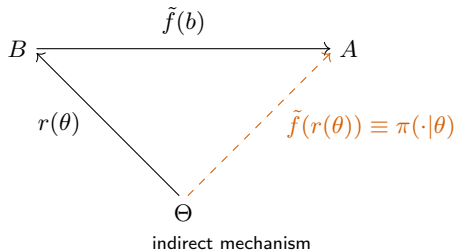
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- Many real world mechanisms are not truthful (e.g., first price auctions)
- not clear that truthful mechanisms are *better* (e.g., second price auctions) (c.f., Li, 2017, Akbarpour & Li, 2020)

So why the obsession with the revelation principle?

- Truthful mechanisms are a good first cut abstraction,
- It is a recipe for constructing *algorithms* that implement allocations,
- It transforms an equilibrium problem into a *constrained* optimization problem.
- From the design perspective, if I cannot find a truthful mechanism that implements my desired rule then *no* mechanism does.



- Sponsored search auctions
- display advertising
- FCC spectrum auctions
- Kidney exchange
- Healthcare systems
- Recommendation systems
- Routing on the Internet
- Resource allocation in the cloud
- Platform design for a sharing economy
- Energy and electricity markets
- Bitcoin
- Participatory democracy
- Crowdsourcing

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- what the designer learns today, they can use tomorrow: **ratchet effect**
 - *e.g., forward-looking bidders understand that bids today determine reserve prices tomorrow ⇒ additional incentive to shave bids above and beyond the strategic and dynamic interaction*

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 - *e.g., sale of a durable good*

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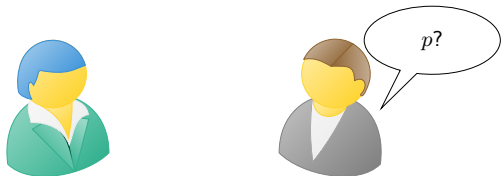
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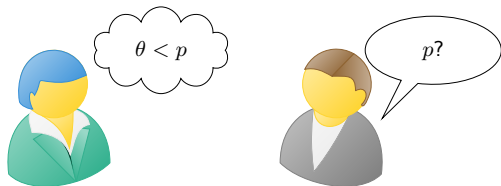


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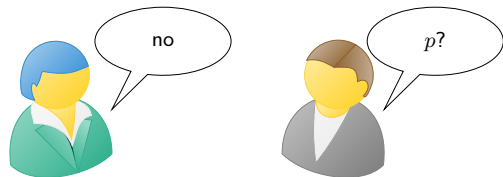


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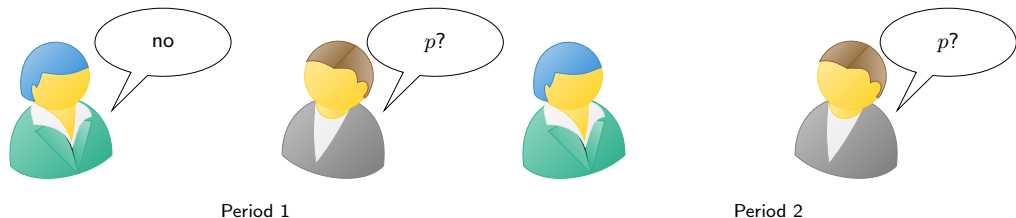


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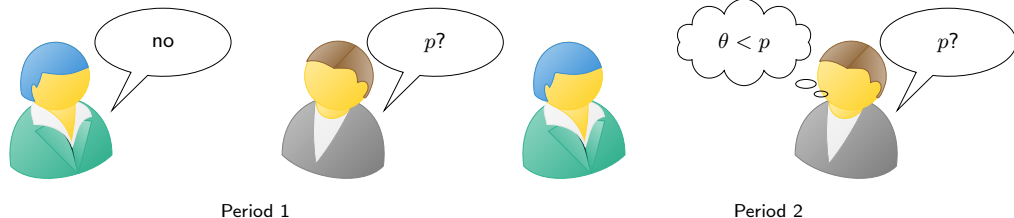
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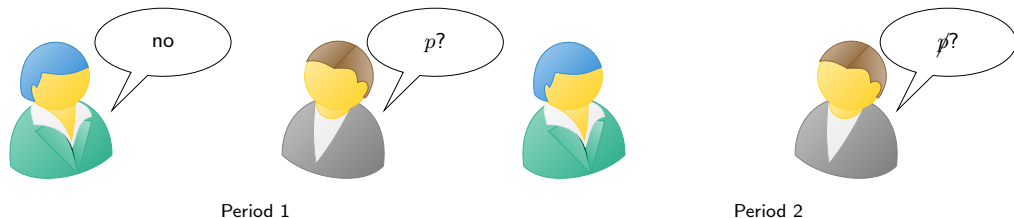
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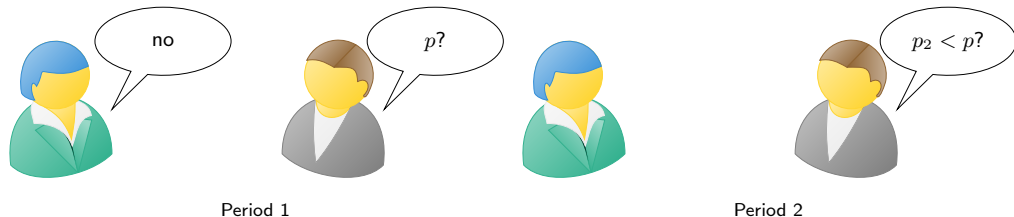
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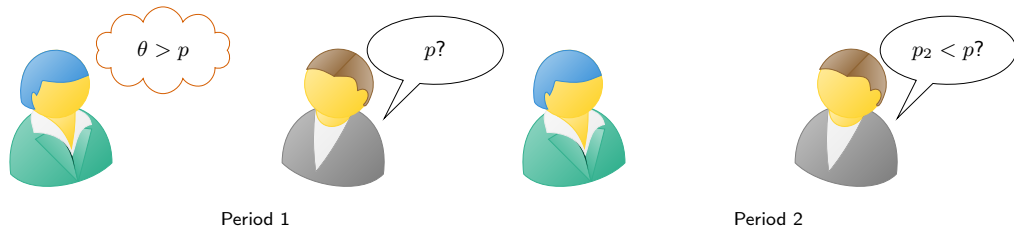
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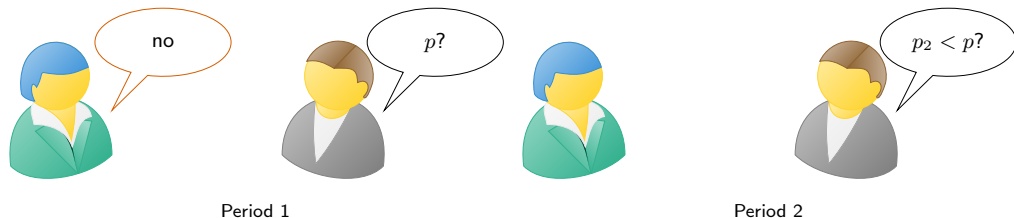
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Oftentimes, optimal mechanisms are not *sequentially rational*:

- i.e., if we gave the designer the possibility to revise the decision rule given the new learned information, they would have an incentive to do so.

There are many examples with these features

- dynamic (ad) auctions (e.g., Google) (c.f., Kanoria & Nazerzadeh, 2014; Papadimitrou et al, 2014; Balseiro et al, 2022)
- repeated sales (e.g., Lobel & Paes Leme, 2017; Devanur et al, 2019; Immorlica et al, 2017)
- procurement (e.g., Gur et al, 2022)

Desiderata: a theory of mechanism design that does not rely so strongly on the assumption that the designer has full commitment

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- Full-term mechanisms w/ 1-sided renegotiation (e.g., Baron & Besanko, 1987)
- Long, but not full, term contract w/ renegotiation (e.g., Rey & Salanie, 1990)
- Cannot commit even to today's mechanism (e.g., Adams & Schwarz 2007, Vartianen 2013, Akbarpour & Li 2020)

Papers in CS & OR that study dynamic lack of commitment focus on this case as well:
Papadimitrou et al, 2014; Lobel & Paes Leme, 2017; Devanur et al, 2019; Immorlica et al, 2017;
Balseiro et al, 2022; Gur et al, 2022

Setting:

- Uninformed designer interacts with privately & persistently informed agent over time
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Examples:

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2. Procurement
3. Political Economy; e.g., taxation and social insurance,
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Few papers analyze optimal mechanisms under limited commitment:

- **Optimal mechanisms w/ finite horizon**, e.g.,
 - Kumar (1985), Laffont & Tirole (1988), Bester and Strausz (2000,2001,2007), Hart & Tirole (1988), Skreta (2006,2015), Bisin & Rampini (2006), Deb & Said (2015), Fiocco & Strausz (2015), Beccutti & Möller (2018)
- **Infinite Horizon under restrictions**, e.g.,
 - Acharya & Ortner (2017), Gerardi & Maestri (2018)
 - **iid private information**: e.g., Sleet and Yeltekin (2006, 2008), Farhi, Sleet, Yeltekin, and Werning (2012), Golosov and Iovino (2021)

The second issue is that the revelation principle no longer holds under limited commitment:

The lack of commitment in repeated adverse-selection situations leads to substantial difficulties for contract theory.

Laffont & Tirole, 1993

- Substantial setback in terms of what we know about optimal policies under [limited commitment](#).

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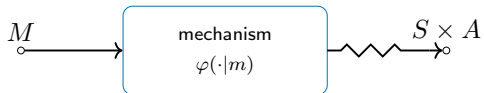
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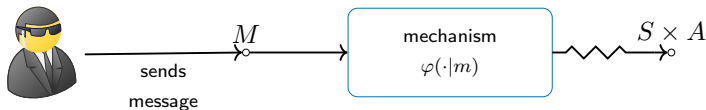
Revelation principle for mechanism design with limited commitment

We characterize a class of mechanisms and strategies that are enough to implement any outcome distribution that can be implemented under limited commitment.



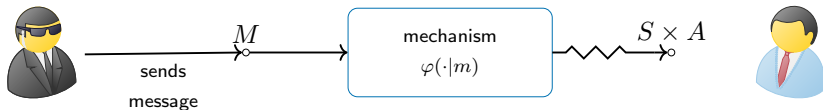
Mechanisms (Myerson '82, Forges '85)

- M is a set of input messages,
- S is a set of output messages,
- φ assigns to each input message a *joint* distribution over output messages and allocations



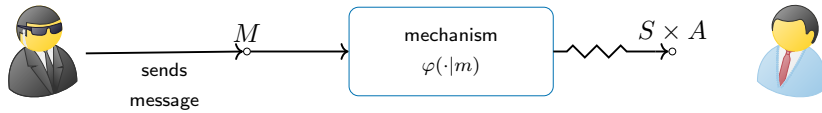
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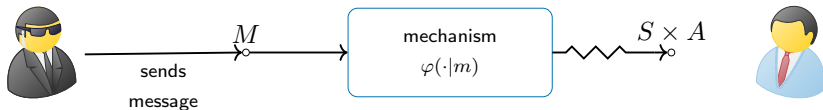
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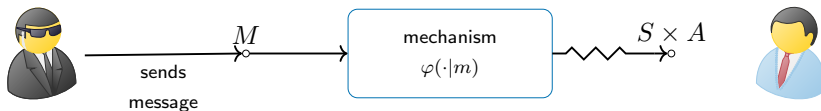




Revelation Principle under commitment

Without loss of generality,

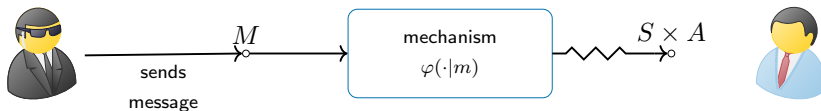
- Communication is **direct**, i.e., $M = \Theta$.
- Communication is **observable**: M and S have the same cardinality and φ is *invertible*.
- Equilibrium communication is **truthful**.



Limited Commitment 1: Bester & Strausz (ECMA, 2001)

Assume:

- Communication is **observable**: M and S have the same cardinality and φ is *invertible*,
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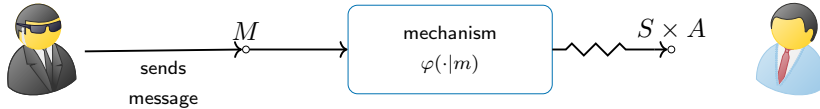
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Then, if the principal earns his highest payoff consistent with the agent's payoff, wlog

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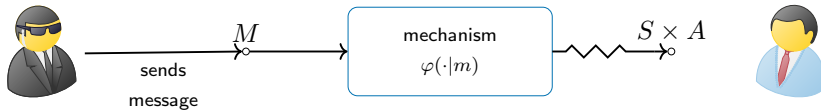
However, ~~Equilibrium communication is truthful~~. (c.f., Papadimitrou et al, 2014)



Limited Commitment 2: Bester & Strausz (JET, 2007)

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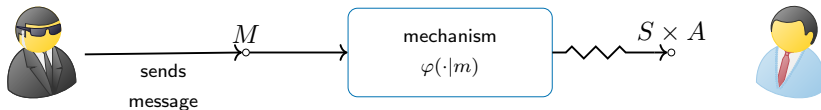
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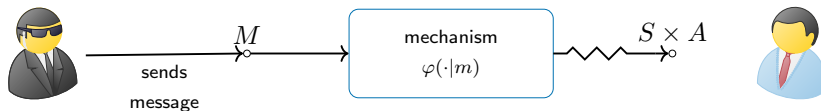
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Revelation Principle for Limited Commitment

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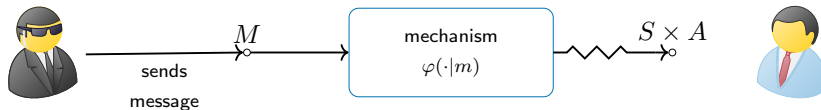
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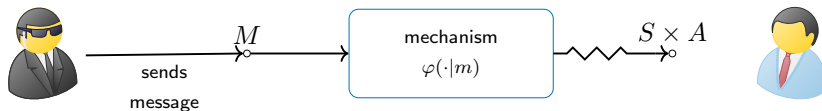
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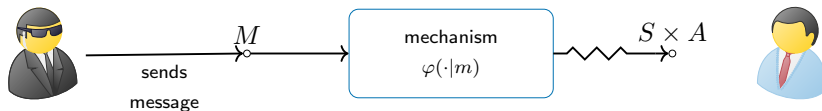


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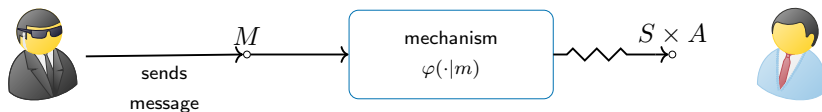


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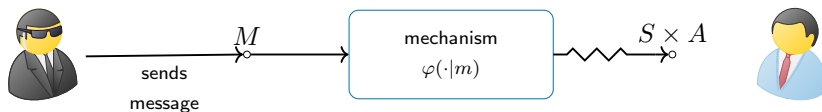


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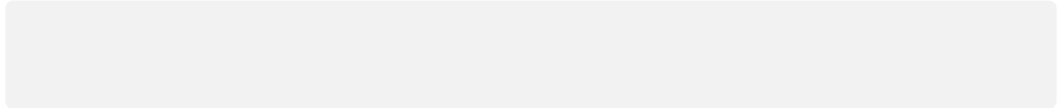
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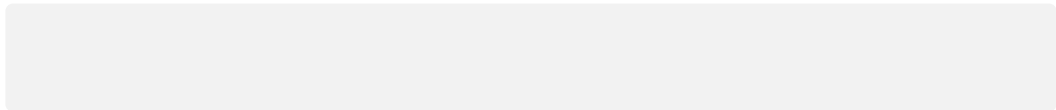
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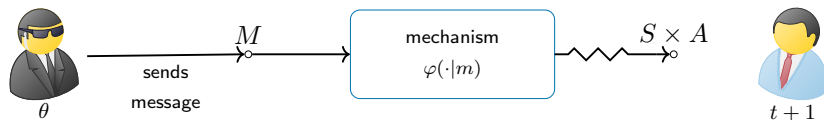
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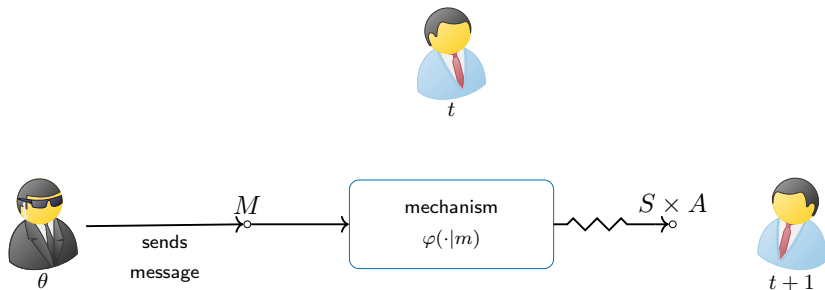
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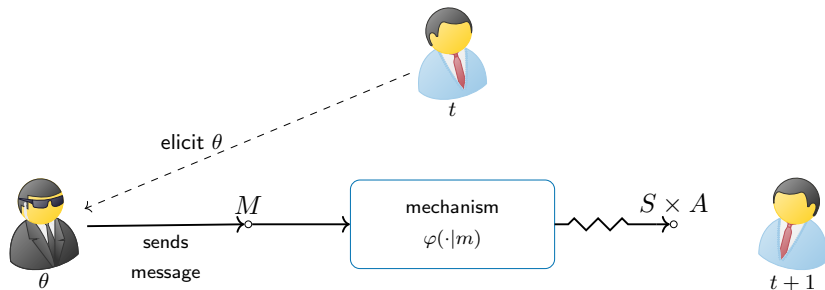
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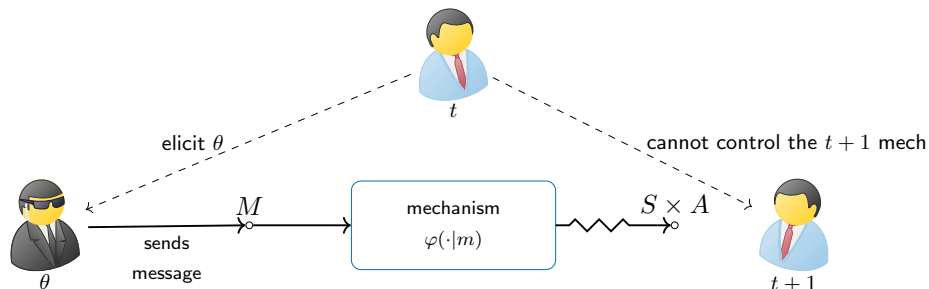
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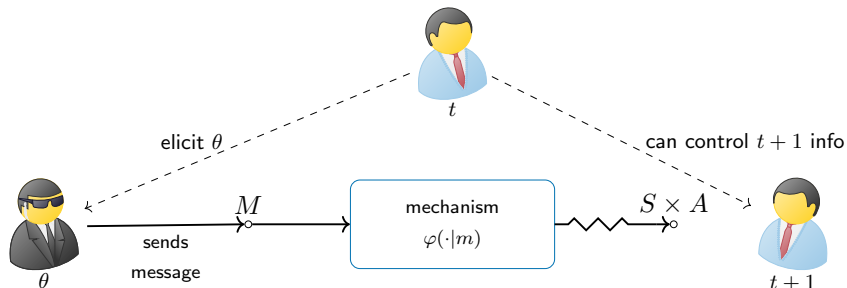
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Revelation Principle for Limited Commitment (Doval & Skreta, 2021)

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 - Truthtelling + participation + **Bayes' plausibility constraint (designer's sequential rationality)**

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- **Today:** Revisit the sale of a durable good w/ a continuum of types and finite horizon

Two final remarks

Two other reasons to care about MDLC in the context of DDDP and AGT:

1. Simplicity
2. Learning

- Limiting the principal's commitment was also an attempt to justify simple mechanisms,
- ... the idea being that it would force the principal to condition his mechanism on less variables (e.g., non-clairvoyant mechanisms, Balseiro et al, 2022)
- It turns out that the optimal mechanism is not necessarily "simpler"
 - e.g., posted prices may no longer be optimal to sell durable goods in finite horizon settings,

- Platforms use **learning algorithms** to optimize on prices/reserve prices based on historical data (c.f., Kanoria & Nazerzadeh, 2014; Haghtalab, Lykouris, Nietert, & Wei, 2022)

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- Our result provides a way of representing these *Bayesian* algorithms and the outcomes that can arise from the strategic interaction with a forward looking agent.
- The analyst is forced to jointly describe the way information is stored and how it is used to determine the allocation.

Sale of a durable good: binary types and two periods

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- The buyer has private information indexed by $\theta \in \Theta \equiv \{\theta_L, \theta_H\}$ and $\mu_0 = \Pr(\theta = \theta_H)$

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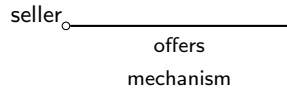
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- If the final allocation is $\{(q_t, x_t)\}_{t \in \{1, 2\}}$, buyer and seller's payoffs are

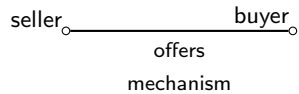
$$U(\cdot, \theta) = \sum_{t=1}^2 \delta^{t-1} (q_t \theta - x_t) \quad \text{and} \quad W(\cdot, \theta) = \sum_{t=1}^2 \delta^{t-1} x_t$$

where $\delta \in (0, 1)$ is a common discount factor.

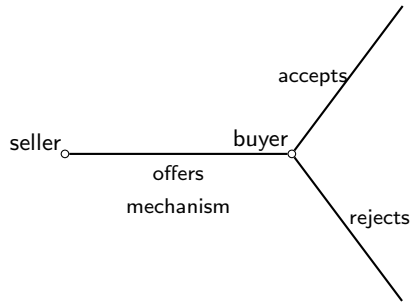
Timing: At the beginning of each period $t \in \{1, 2\}$



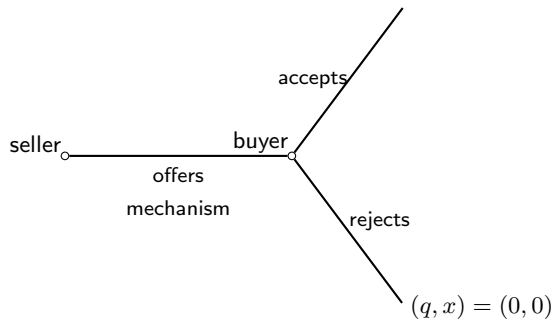
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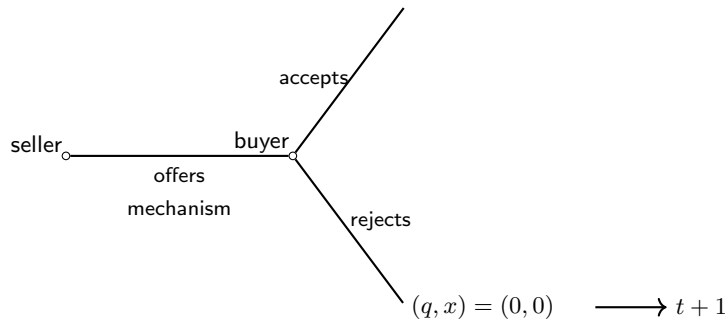
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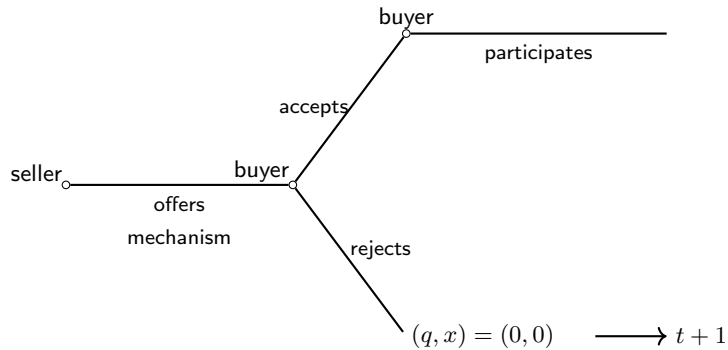
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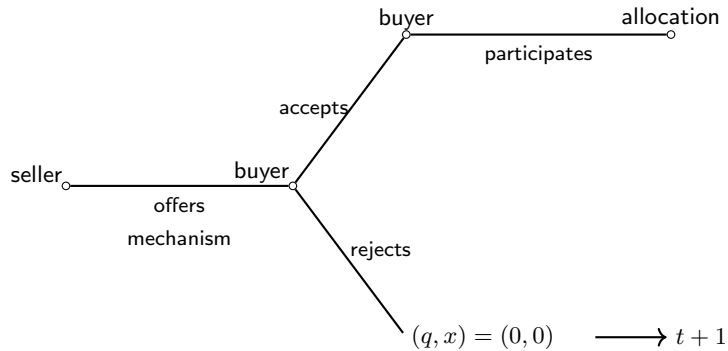
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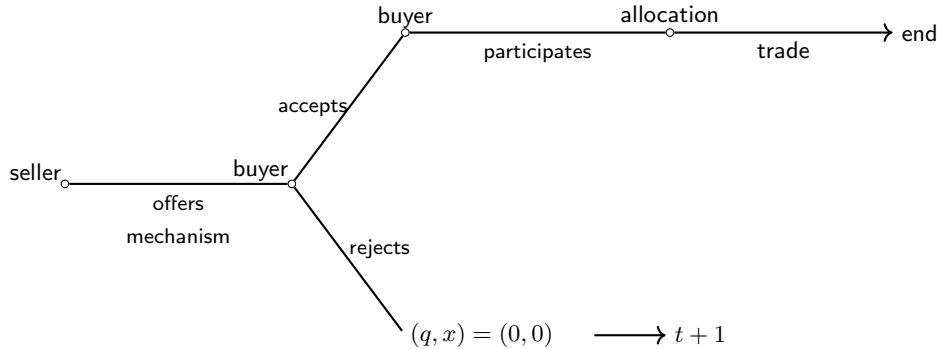
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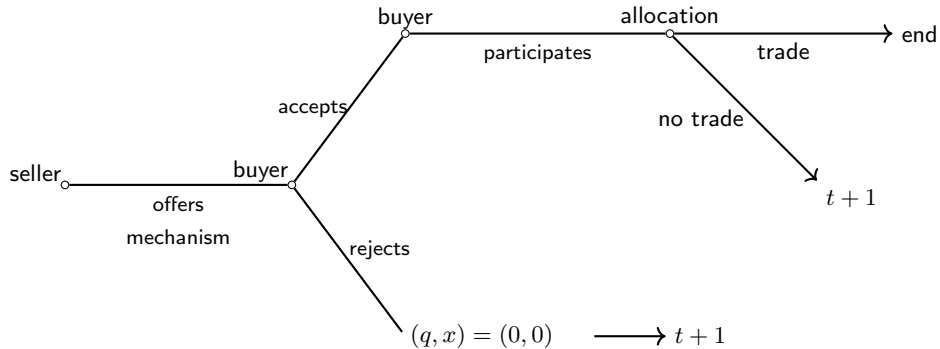
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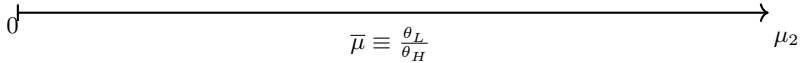


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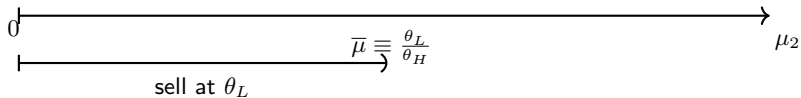


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- Let μ_2 denote the seller's belief that $\theta = \theta_H$.

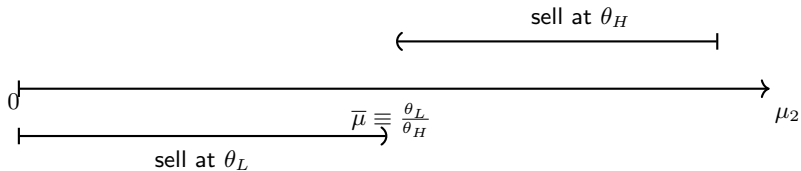
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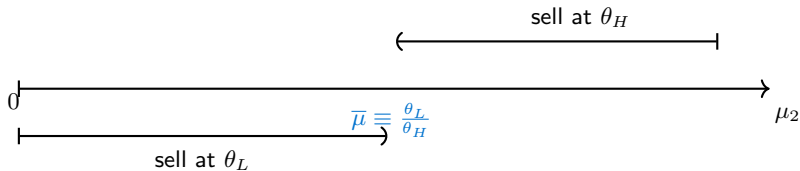
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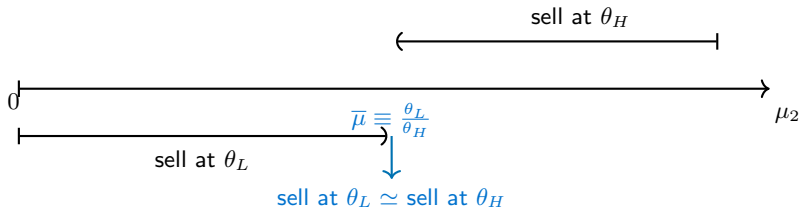
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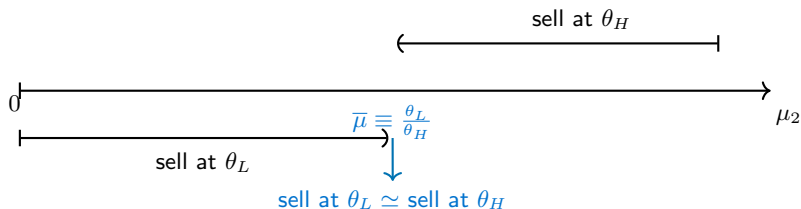
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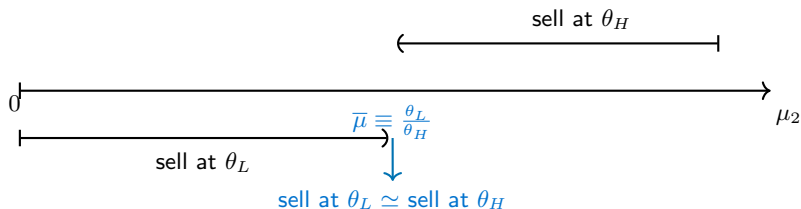
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- Why $\bar{\mu}$? Whenever the seller sells to both types, he leaves rents $\mu_2 \Delta \theta$ to θ_H .

$$\begin{aligned} \theta_L &= \mu_2(\theta_H - \Delta\theta) + (1 - \mu_2)\theta_L = \mu_2\theta_H + (1 - \mu_2)\left(\theta_L - \frac{\mu_2}{1 - \mu_2}\Delta\theta\right) \\ &= \mu_2\theta_H + (1 - \mu_2)\hat{\theta}_L(\mu_2) \end{aligned}$$

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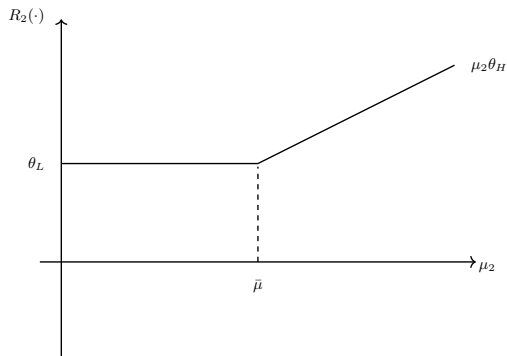
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When $\mu_2 = \bar{\mu}$, then $\hat{\theta}_L(\mu_2) = 0$.

Wrapping up:

$$R_2(\mu_2) = \begin{cases} \theta_L & \text{if } \mu_2 \leq \bar{\mu} \\ \mu_2\theta_H & \text{if } \mu_2 > \bar{\mu} \end{cases} = \begin{cases} \mu_2\theta_H + (1 - \mu_2)\hat{\theta}_L(\mu_2) & \text{if } \mu_2 \leq \bar{\mu} \\ \mu_2\theta_H & \text{if } \mu_2 > \bar{\mu} \end{cases}$$



Seller's payoff in period 2

- Recall μ_1 is the prior probability that $\theta = \theta_H$.

▶ $t = 2$

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- A **mechanism** is a tuple

$$M \xrightarrow{\varphi(\cdot|m)} S \times A$$

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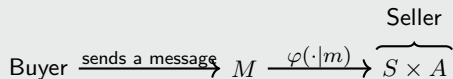
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$$\text{Buyer} \xrightarrow{\text{sends a message}} M \xrightarrow{\varphi(\cdot|m)} S \times A$$

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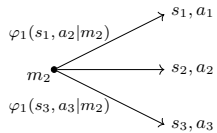
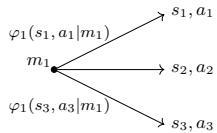
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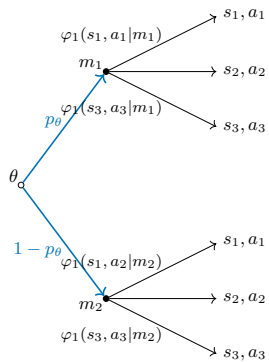
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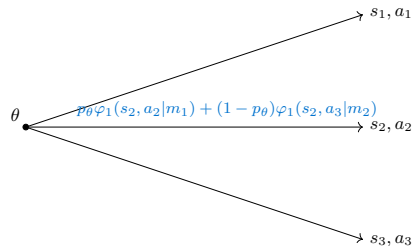
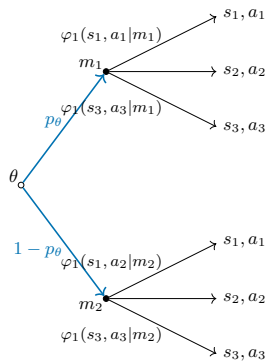
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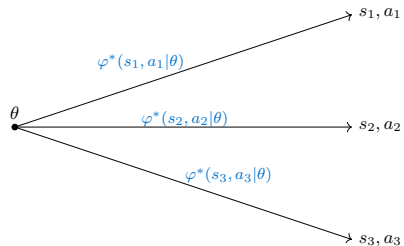
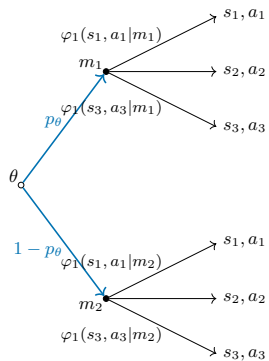
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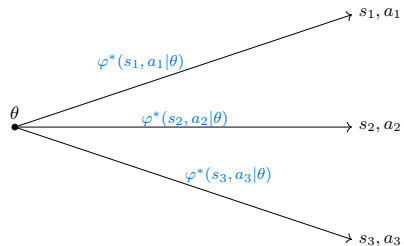
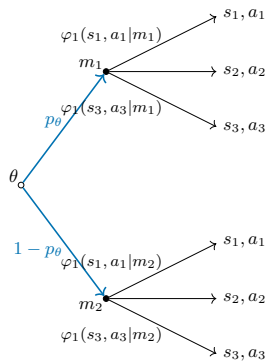
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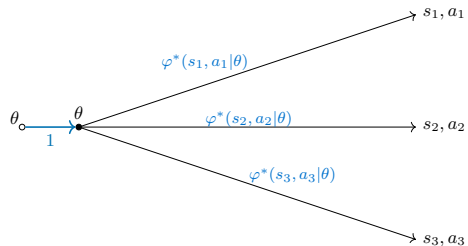
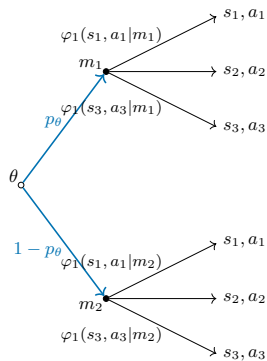
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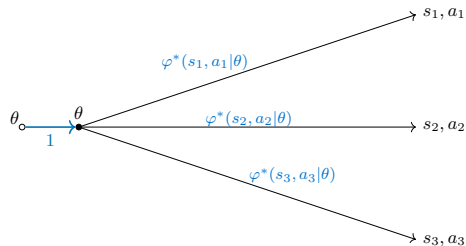
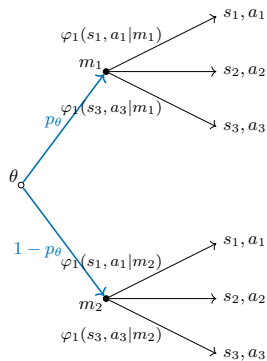
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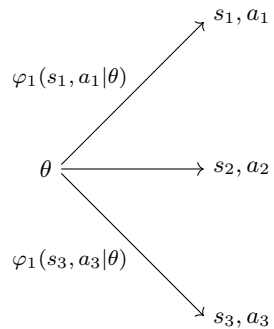
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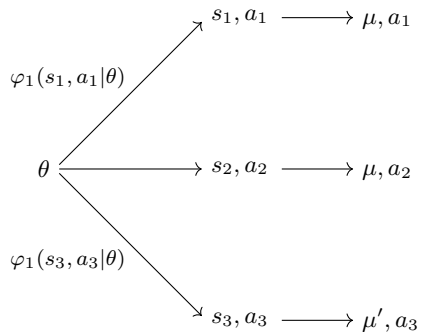


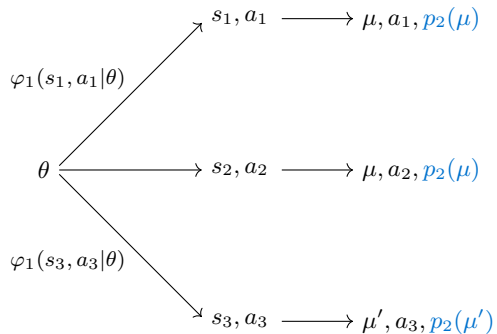
- Bester and Strausz (JET, 2007)

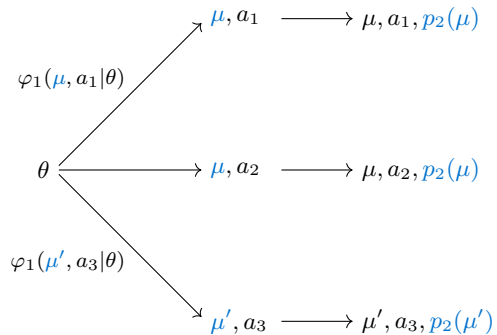
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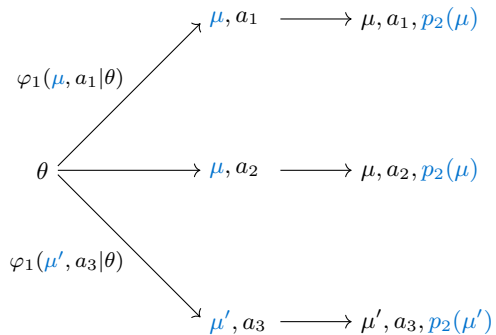


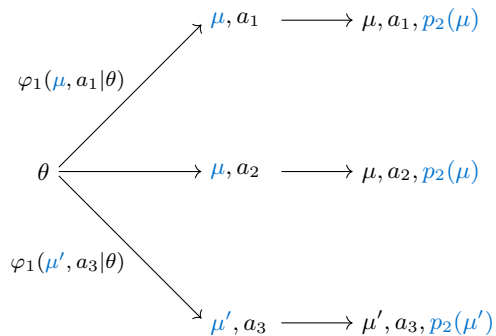
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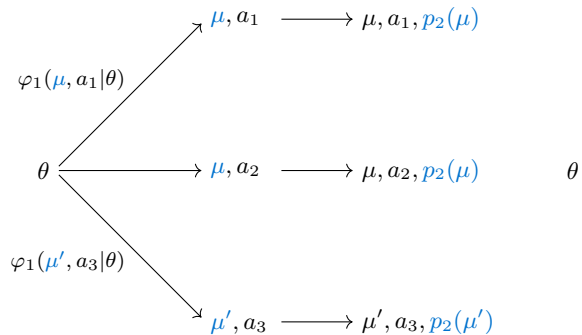
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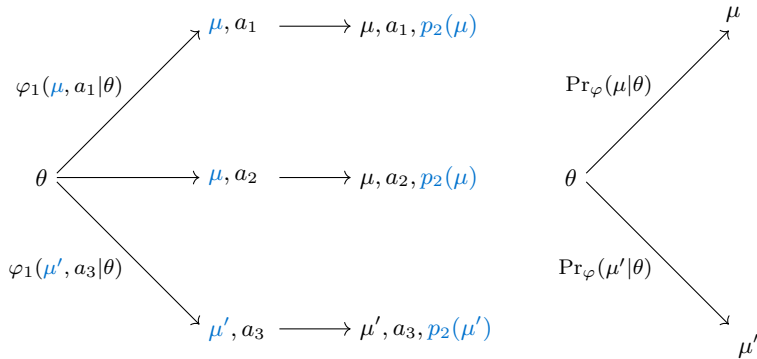
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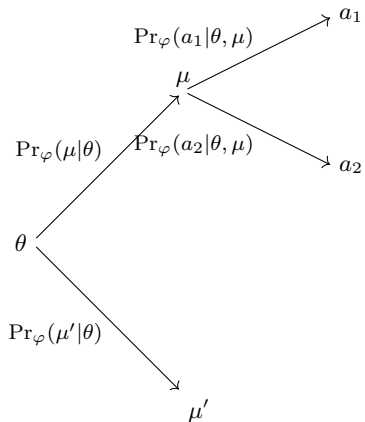
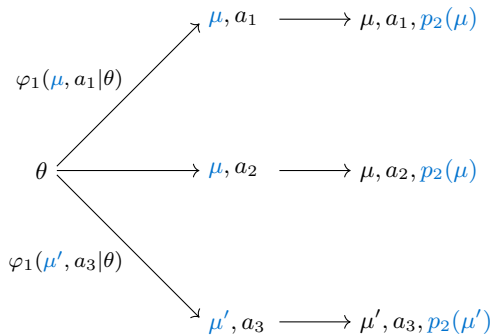




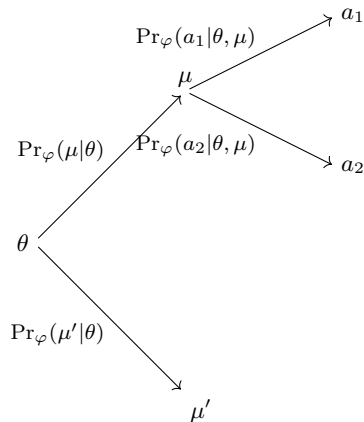
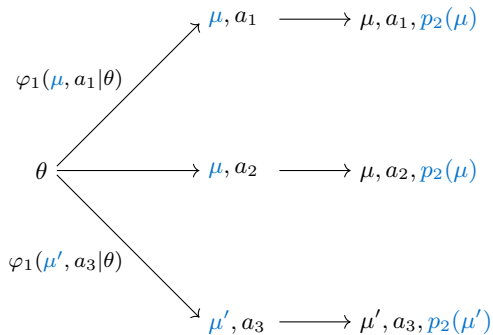




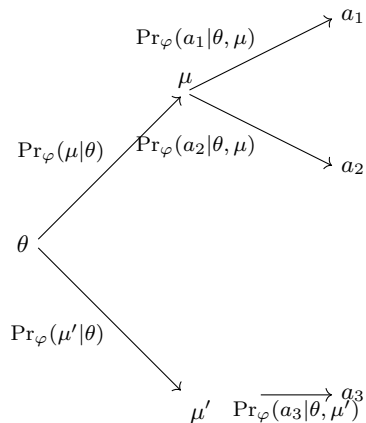
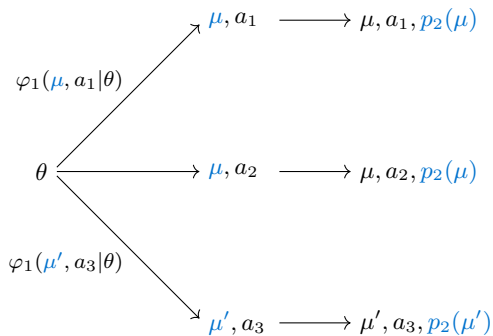
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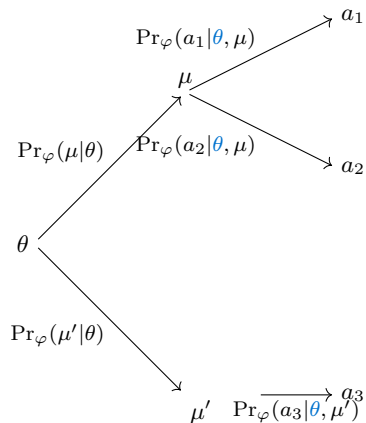
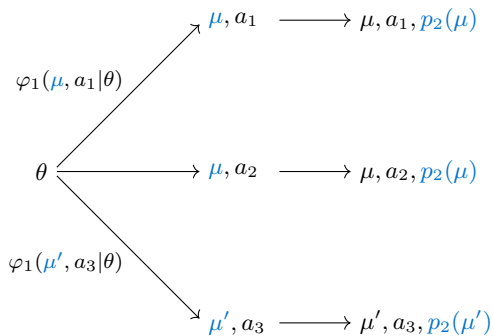
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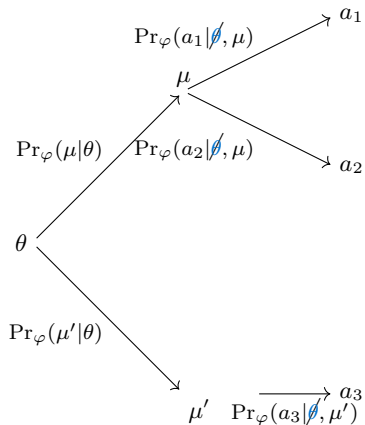
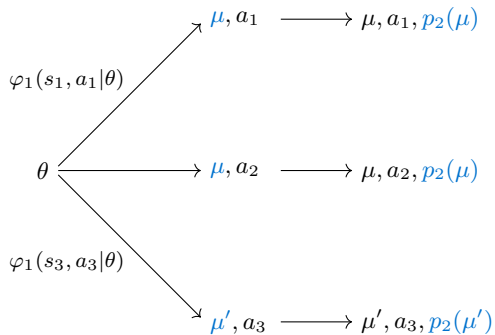
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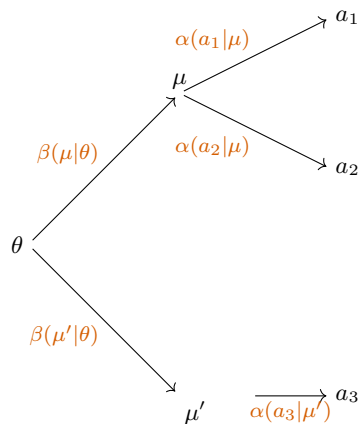
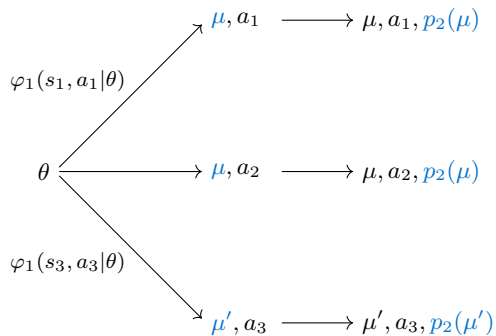
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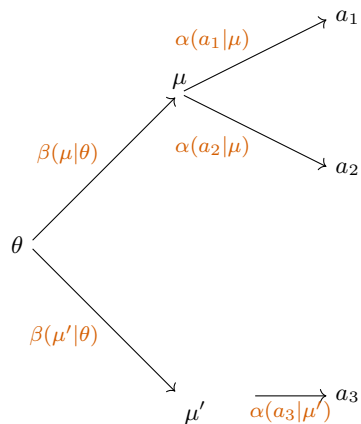
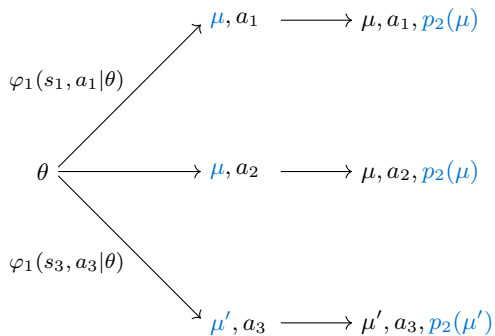
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- Separate the design of the information from that of the allocation
- β is the mechanism's **disclosure rule** and α is the mechanism's **allocation rule**.

Quasilinearity + separation between allocation and information:

- No need to randomize on transfers: $x(\mu_2)$ is the (expected) payment when output message is μ_2
- $q(\mu_2)$ is the probability of selling the good when output message is μ_2

Thus, the seller's optimal outcome solves:

$$\max_{\text{mechanisms}} \text{Revenue}$$

where $M_1 = \Theta$, $S_1 = \Delta(\Theta)$, $\varphi = \beta \otimes \alpha$ subject to

- Participation
- Truthtelling
- Consistency between beliefs and output messages.

Seller optimal outcome

constrained optimization

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$$\left(\sum_{\theta \in \Theta} \mu_1(\theta) \beta(\mu_2 | \theta) \right) [x(\mu_2) + (1 - q(\mu_2)) \delta R_2(\mu_2)]$$

Thus, the seller's optimal outcome solves:

$$\max_{\beta, q, x} \sum_{\mu_2 \in \Delta(\Theta)} \left(\sum_{\theta \in \Theta} \mu_1(\theta) \beta(\mu_2 | \theta) \right) [x(\mu_2) + (1 - q(\mu_2)) \delta R_2(\mu_2)]$$

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Participation $_{\theta}$:
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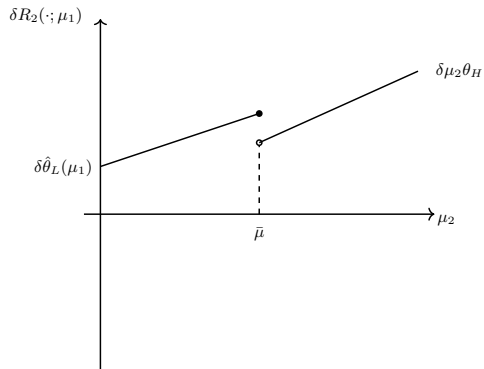
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Sale of a durable good: $t = 1$

$$\delta R_2(\mu_2; \mu_1) = \begin{cases} \delta(\mu_2\theta_H + (1 - \mu_2)\hat{\theta}_L(\mu_1)) & \text{if } \mu_2 < \bar{\mu} \\ \delta\mu_2\theta_H & \text{if } \mu_2 > \bar{\mu} \end{cases}$$

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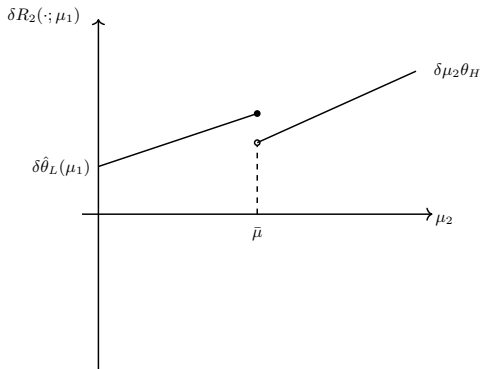
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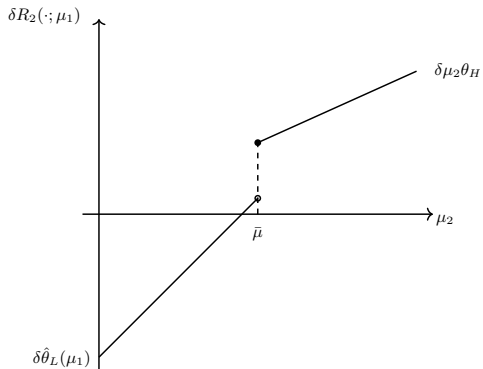
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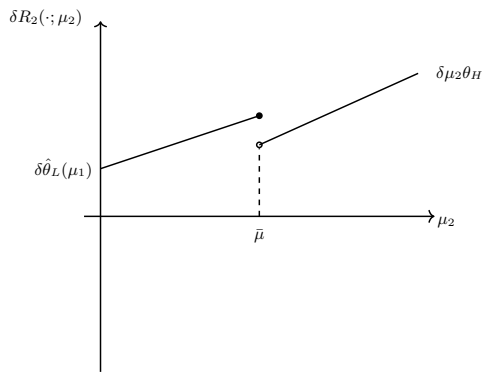
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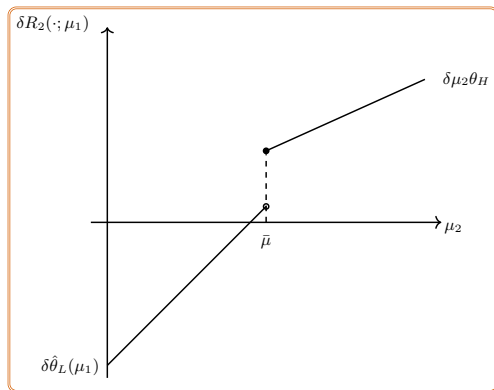
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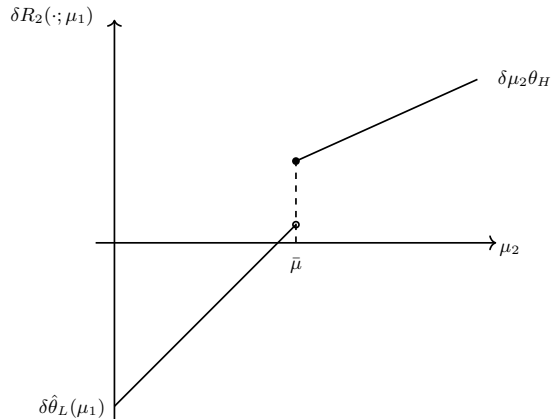
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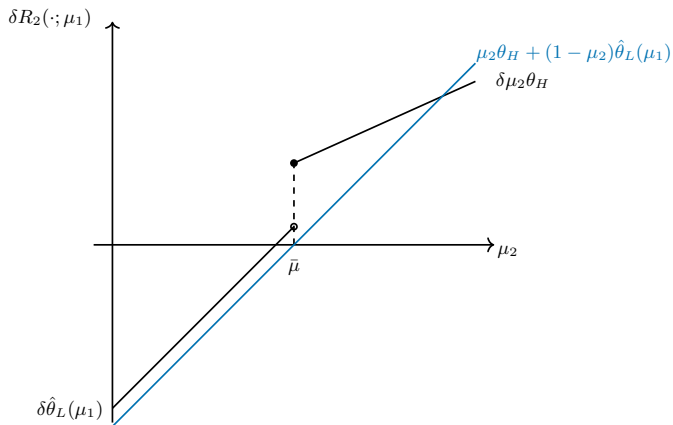
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$$\max_{\tau, q} \sum_{\mu_2 \in \Delta(\Theta)} \tau(\mu_2) \left[q(\mu_2)(\mu_2 \theta_H + (1 - \mu_2) \hat{\theta}_L(\mu_1)) + (1 - q(\mu_2)) \delta R_2(\mu_2; \mu_1) \right]$$

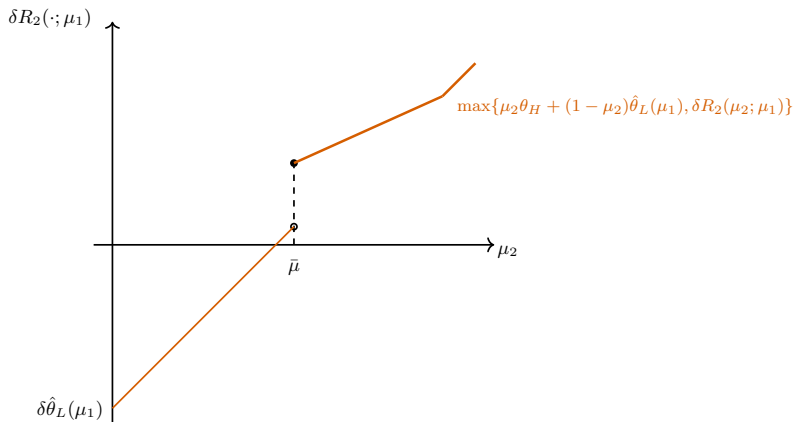
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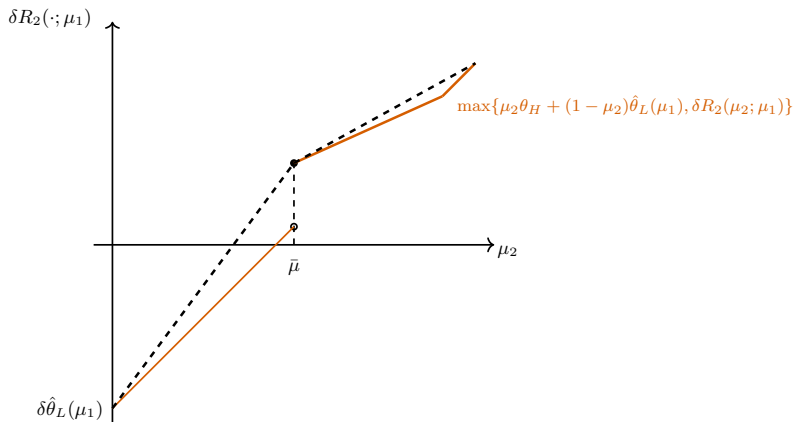
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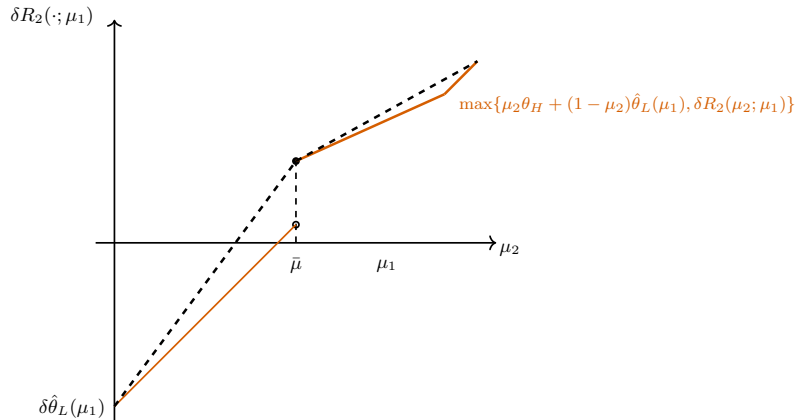
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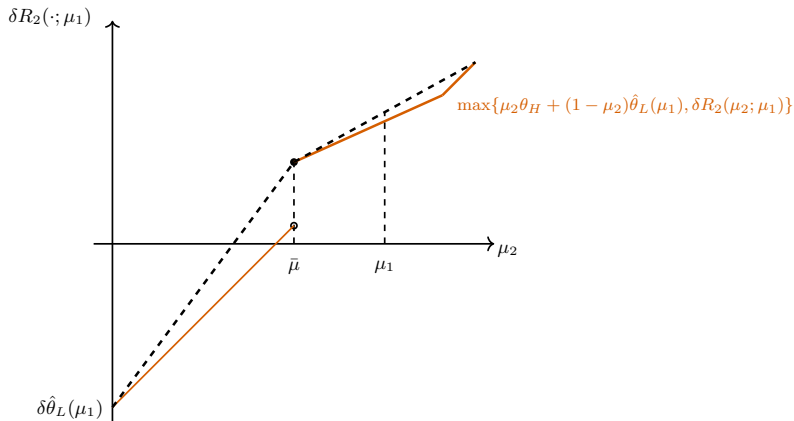
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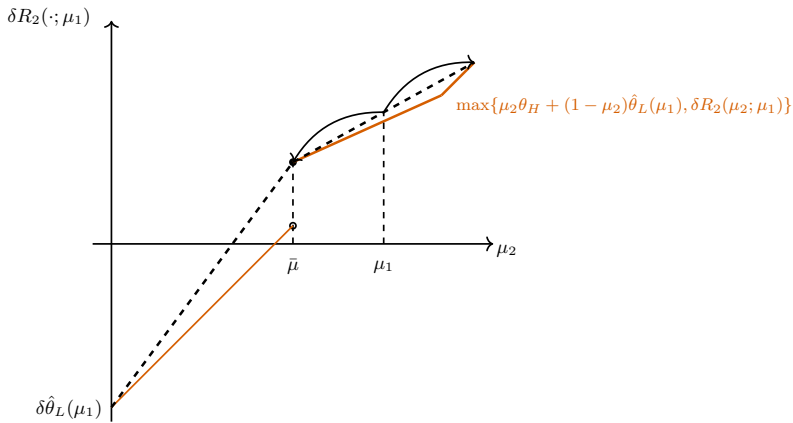
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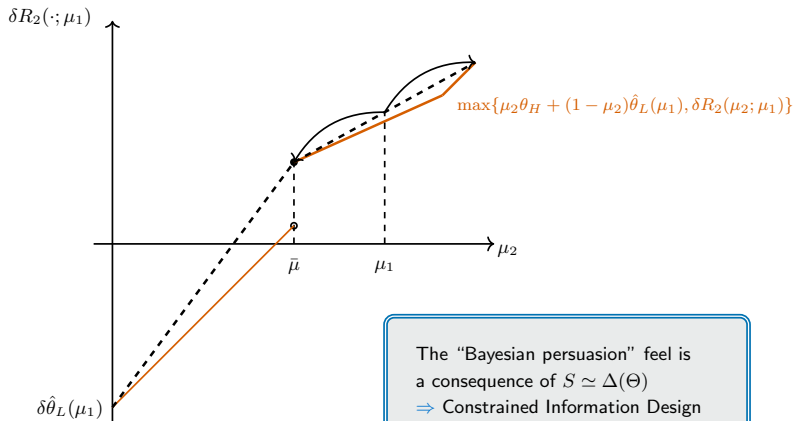
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- Seller splits μ_1 between $\mu_2 = \bar{\mu}$ and $\mu_2 = 1$
- He sells when $\mu_2 = 1$ ($q(1) = 1$) and delays when $\mu_2 = \bar{\mu}$ ($q(\bar{\mu}) = 0$)
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The “Bayesian persuasion” feel is
 a consequence of $S \simeq \Delta(\Theta)$
 \Rightarrow Constrained Information Design

Economic trade-off: tailor the allocation to the agent's report vs. learning about the agent's type.

- No such trade-off when there is commitment: acquired information can always be “forgotten.”
- The seller slows down learning:
 - Similar to Kanoria & Nazerzadeh, 2014; Abernethy et al., 2019; Haghtalab, Lykouris, Nietert, & Wei, 2022

Open questions

This is a problem that had been open in Economics for 30 years. There's much to do!

1. Most glaring: multiple agents (the existing counterexamples do not survive with our mechanisms)
 - How to aggregate the information from the multiple agents? (e.g., Halpern & Teague, 2006)
2. More practical: How to implement direct-Blackwell mechanisms?
 - Multiple (infinite?) rounds of indirect observable communication?
 - Cryptographic commitments? (e.g., Ferreira & Weinberg, 2020)