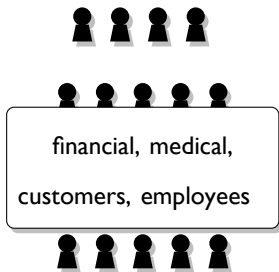


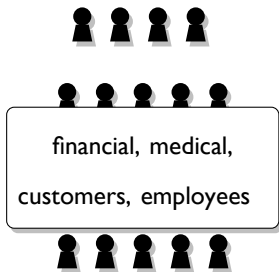
encrypted computation *from* lattices



Hoeteck Wee

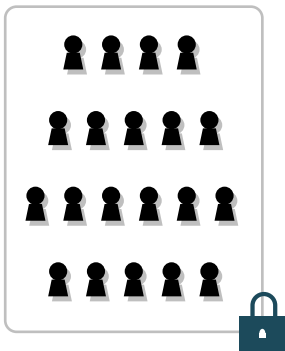


BIG DATA



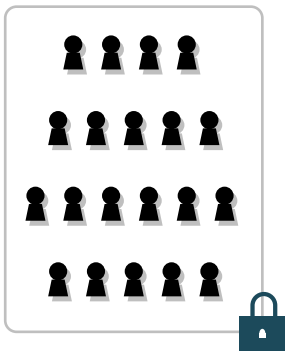
BIG DATA

Q. privacy.



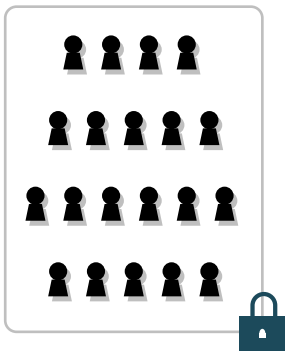
BIG DATA

Q. privacy.



BIG DATA

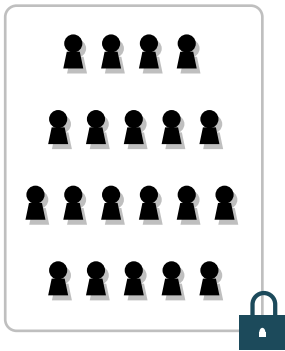
Q. privacy. utility?



BIG DATA

Q. privacy + utility

encrypted computation

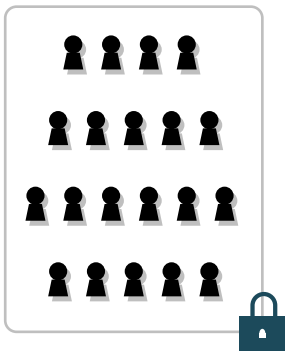


BIG DATA

Q. privacy + utility

encrypted computation

3 notions

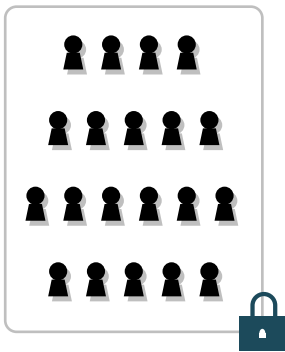


BIG DATA

Q. privacy + utility

encrypted computation

3 notions *from* **lattices**



BIG DATA

Q. privacy + utility

encrypted computation

3 notions + **1** equation

fully **homomorphic** encryption

fully homomorphic **encryption**

syntax. $\text{enc}(\text{sk}, \cdot)$, $\text{dec}(\text{sk}, \cdot)$

functionality.

fully homomorphic **encryption**

syntax. $\text{enc}(\text{sk}, \cdot)$, $\text{dec}(\text{sk}, \cdot)$

functionality. $\text{dec}(\text{sk}, \text{enc}(\text{sk}, x)) = x$

fully homomorphic **encryption**

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{dec}(\text{sk}, \text{enc}(\text{sk}, x)) = x$

fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \xrightarrow{\text{eval}} \text{enc}(\text{sk}, f(x))$

fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \xrightarrow{\text{eval}} \text{enc}(\text{sk}, f(x))$

FHE for **circuits** from lattices

[Gentry 09, Brakerski Vaikuntanathan 11, Gentry Sahai Waters 13]

fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \xrightarrow{\text{eval}} \text{enc}(\text{sk}, f(x))$

FHE for **circuits** from **LWE**

[Gentry 09, Brakerski Vaikuntanathan 11, Gentry Sahai Waters 13]

$$(\mathbf{B}, \mathbf{sB} + \mathbf{e}) \approx_c \text{uniform}$$



fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \xrightarrow{\text{eval}} \text{enc}(\text{sk}, f(x))$

FHE for **circuits** from **LWE**

[Gentry 09, Brakerski Vaikuntanathan 11, Gentry Sahai Waters 13]

fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

1

fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

1

t
 sk

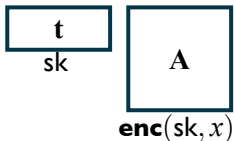
over \mathbb{Z}_q

fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

1



fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

1

$$\begin{array}{c} \boxed{\mathbf{t}} \\ \text{sk} \end{array} \begin{array}{c} \boxed{\mathbf{A}} \\ \text{enc}(\text{sk}, x) \end{array} = \begin{array}{c} \boxed{x \mathbf{t}} \\ \mathbf{t}: \text{eigenvector} \end{array}$$

fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

1

$$\begin{array}{c} \boxed{\mathbf{t}} \\ \text{sk} \end{array} \quad \boxed{\mathbf{A}_i} \quad = \quad \boxed{x_i \mathbf{t}}$$

$\text{enc}(\text{sk}, x_i)$ \mathbf{t} : eigenvector

$$\text{enc}(\text{sk}, x_1), \text{enc}(\text{sk}, x_2) \stackrel{?}{\mapsto} \text{enc}(\text{sk}, x_1 + x_2), \text{enc}(\text{sk}, x_1 x_2)$$

fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

1

$$\begin{array}{c} \boxed{\mathbf{t}} \\ \text{sk} \end{array} \cdot \begin{array}{c} \boxed{\mathbf{A}_i} \\ \text{enc}(\text{sk}, x_i) \end{array} = \boxed{x_i \mathbf{t}}$$

addition: $\mathbf{t} \cdot (\mathbf{A}_1 + \mathbf{A}_2) = (x_1 + x_2)\mathbf{t}$

fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

1

$$\begin{array}{c} \boxed{\mathbf{t}} \\ \text{sk} \end{array} \cdot \begin{array}{c} \boxed{\mathbf{A}_i} \\ \text{enc}(\text{sk}, x_i) \end{array} = \boxed{x_i \mathbf{t}}$$

addition: $\mathbf{t} \cdot (\mathbf{A}_1 + \mathbf{A}_2) = (x_1 + x_2)\mathbf{t}$

multiplication: $\mathbf{t} \cdot \quad = x_1 x_2 \mathbf{t}$

fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

1

$$\begin{array}{c} \boxed{\mathbf{t}} \\ \text{sk} \end{array} \cdot \begin{array}{c} \boxed{\mathbf{A}_i} \\ \text{enc}(\text{sk}, x_i) \end{array} = \boxed{x_i \mathbf{t}}$$

addition: $\mathbf{t} \cdot (\mathbf{A}_1 + \mathbf{A}_2) = (x_1 + x_2)\mathbf{t}$

multiplication: $\mathbf{t} \cdot \mathbf{A}_1 \mathbf{A}_2 = x_1 x_2 \mathbf{t}$

$$\text{LHS} = x_1 \mathbf{t} \cdot \mathbf{A}_2 = \dots$$

fully homomorphic encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

1

$$\begin{array}{c} \boxed{\mathbf{t}} \\ \text{sk} \end{array} \cdot \begin{array}{c} \boxed{\mathbf{A}_i} \\ \text{enc}(\text{sk}, x_i) \end{array} = \boxed{x_i \mathbf{t}}$$

addition: $\mathbf{t} \cdot (\mathbf{A}_1 + \mathbf{A}_2) = (x_1 + x_2)\mathbf{t}$

multiplication: $\mathbf{t} \cdot \mathbf{A}_1\mathbf{A}_2 = x_1x_2\mathbf{t}$

polynomials: $\mathbf{t} \cdot (\mathbf{A}_1\mathbf{A}_2 + \mathbf{A}_3\mathbf{A}_4) = (x_1x_2 + x_3x_4)\mathbf{t}$

fully homomorphic encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

1

$$\boxed{\begin{array}{c} \mathbf{t} \\ \text{sk} \end{array}} \cdot \boxed{\begin{array}{c} \mathbf{A}_i \\ \text{enc}(\text{sk}, x_i) \end{array}} = \boxed{x_i \mathbf{t}}$$

addition: $\mathbf{t} \cdot (\mathbf{A}_1 + \mathbf{A}_2) = (x_1 + x_2)\mathbf{t}$

multiplication: $\mathbf{t} \cdot \mathbf{A}_1 \mathbf{A}_2 = x_1 x_2 \mathbf{t}$

polynomials: $\mathbf{t} \cdot \underbrace{f(\mathbf{A}_1, \dots, \mathbf{A}_n)}_{\mathbf{A}_f} = f(x_1, \dots, x_n)\mathbf{t}$

fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

2

+ noise

$$\begin{array}{ccc} \boxed{\begin{array}{c} \mathbf{t} \\ \text{sk} \end{array}} & \boxed{\begin{array}{c} \mathbf{A}_i \\ \text{enc}(\text{sk}, x_i) \end{array}} & \approx & \boxed{\begin{array}{c} x_i \mathbf{t} \end{array}} \end{array}$$

addition: $\mathbf{t} \cdot (\mathbf{A}_1 + \mathbf{A}_2) = (x_1 + x_2)\mathbf{t}$

multiplication: $\mathbf{t} \cdot \mathbf{A}_1\mathbf{A}_2 = x_1x_2\mathbf{t}$

fully **homomorphic** encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

2

+ noise

$$\begin{array}{ccc} \boxed{\begin{array}{c} \mathbf{t} \\ \text{sk} \end{array}} & \boxed{\begin{array}{c} \mathbf{A}_i \\ \text{enc}(\text{sk}, x_i) \end{array}} & \approx & \boxed{\begin{array}{c} x_i \mathbf{t} \end{array}} \end{array}$$

addition: $\mathbf{t} \cdot (\mathbf{A}_1 + \mathbf{A}_2) \approx (x_1 + x_2)\mathbf{t}$

– *proof.* small + small = small

fully homomorphic encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

2

+ noise

$$\begin{array}{ccc} \boxed{\begin{array}{c} \mathbf{t} \\ \text{sk} \end{array}} & \boxed{\begin{array}{c} \mathbf{A}_i \\ \text{enc}(\text{sk}, x_i) \end{array}} & \approx & \boxed{x_i \mathbf{t}} \end{array}$$

addition: $\mathbf{t} \cdot (\mathbf{A}_1 + \mathbf{A}_2) \approx (x_1 + x_2)\mathbf{t}$

multiplication: $\mathbf{t} \cdot \mathbf{A}_1 \mathbf{A}_2 \approx x_1 x_2 \mathbf{t}$

– *proof.* small $\cdot \mathbf{A}_2 = \text{big}$

fully homomorphic encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

3 \mathbf{A}_i small

$$\begin{array}{ccc} \boxed{\begin{array}{c} \mathbf{t} \\ \text{sk} \end{array}} & \boxed{\begin{array}{c} \mathbf{A}_i \\ \text{enc}(\text{sk}, x_i) \end{array}} & \approx & \boxed{x_i \mathbf{t}} \end{array}$$

addition: $\mathbf{t} \cdot (\mathbf{A}_1 + \mathbf{A}_2) \approx (x_1 + x_2)\mathbf{t}$

multiplication: $\mathbf{t} \cdot \mathbf{A}_1\mathbf{A}_2 \not\approx x_1x_2\mathbf{t}$

– *proof.* small $\cdot \mathbf{A}_2 =$ big

fully homomorphic encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

3 \mathbf{A}_i small

$$\begin{array}{ccc} \boxed{\begin{array}{c} \mathbf{t} \\ \text{sk} \end{array}} & \boxed{\begin{array}{c} \mathbf{A}_i \\ \text{enc}(\text{sk}, x_i) \end{array}} & \approx & \boxed{\begin{array}{c} x_i \mathbf{t} \end{array}} \end{array}$$

addition: $\mathbf{t} \cdot (\mathbf{A}_1 + \mathbf{A}_2) \approx (x_1 + x_2)\mathbf{t}$

multiplication: $\mathbf{t} \cdot \mathbf{A}_1 \mathbf{A}_2 \approx x_1 x_2 \mathbf{t}$

– *proof.* $\text{small} \cdot \mathbf{A}_2 = \text{small}$

fully homomorphic encryption

security. $\text{enc}(\text{sk}, x)$ hides x

functionality. $\text{enc}(\text{sk}, x) \mapsto \text{enc}(\text{sk}, f(x))$

3 \mathbf{A}_i small

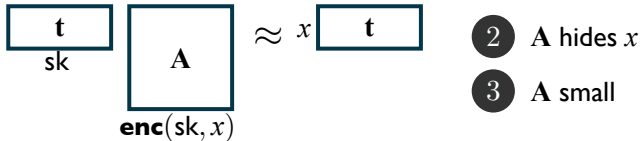
$$\begin{array}{ccc} \boxed{\begin{array}{c} \mathbf{t} \\ \text{sk} \end{array}} & \boxed{\begin{array}{c} \mathbf{A}_i \\ \text{enc}(\text{sk}, x_i) \end{array}} & \approx & \boxed{\begin{array}{c} x_i \mathbf{t} \end{array}} \end{array}$$

addition: $\mathbf{t} \cdot (\mathbf{A}_1 + \mathbf{A}_2) \approx (x_1 + x_2)\mathbf{t}$

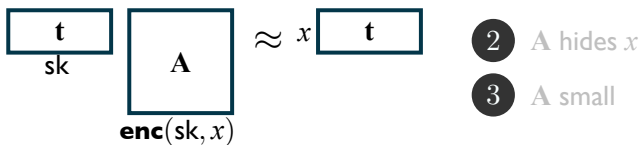
multiplication: $\mathbf{t} \cdot \mathbf{A}_1 \mathbf{A}_2 \approx x_1 x_2 \mathbf{t}$

polynomials: $\mathbf{t} \cdot \underbrace{f(\mathbf{A}_1, \dots, \mathbf{A}_n)}_{\mathbf{A}_f} \approx f(x_1, \dots, x_n)\mathbf{t}$

fully **homomorphic** encryption

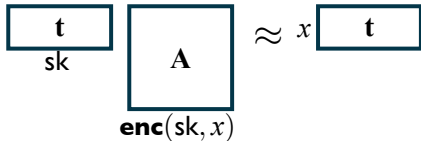


fully **homomorphic** encryption



$$\underbrace{(s - 1)}_t$$

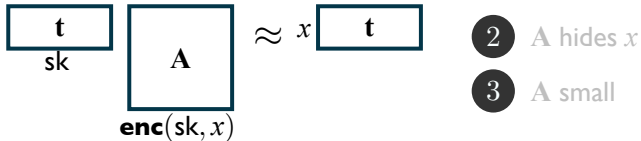
fully **homomorphic** encryption



- 2 A hides x
- 3 A small

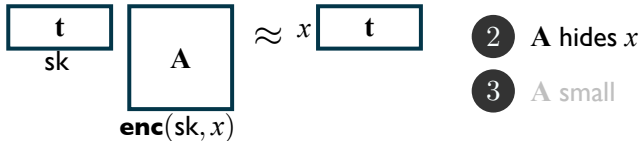
$$\underbrace{\begin{pmatrix} \mathbf{s} & -1 \end{pmatrix}}_t \begin{pmatrix} \mathbf{B} \\ \mathbf{sB} + \mathbf{e} \end{pmatrix} \approx \mathbf{0}$$

fully **homomorphic** encryption



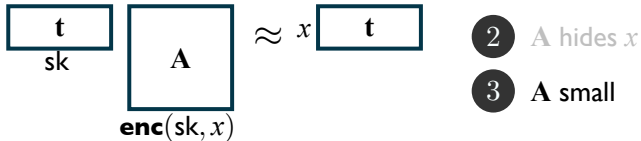
$$\underbrace{\begin{pmatrix} s & -1 \end{pmatrix}}_t \left(\begin{pmatrix} B \\ sB + e \end{pmatrix} + xI \right) \approx xt$$

fully **homomorphic** encryption



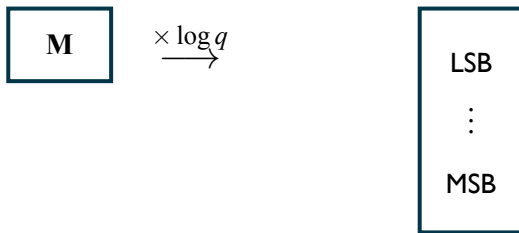
$$\underbrace{(s - 1)}_t \left(\left(\begin{pmatrix} \mathbf{B} \\ s\mathbf{B} + \mathbf{e} \end{pmatrix} + x\mathbf{I} \right) \right) \approx xt$$

fully **homomorphic** encryption

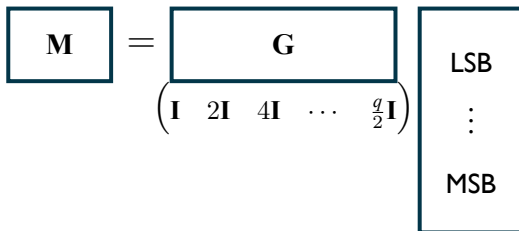


$$\underbrace{(s - 1)}_t \left(\begin{pmatrix} \mathbf{B} \\ s\mathbf{B} + \mathbf{e} \end{pmatrix} + x\mathbf{I} \right) \approx xt$$

fully **homomorphic** encryption



fully **homomorphic** encryption

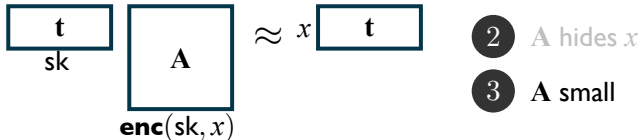


fully **homomorphic** encryption



$$\mathbf{M} = \mathbf{G} \begin{pmatrix} \mathbf{I} & 2\mathbf{I} & 4\mathbf{I} & \cdots & \frac{q}{2}\mathbf{I} \end{pmatrix} \mathbf{G}^{-1}(\mathbf{M})$$

fully **homomorphic** encryption



$$\underbrace{\begin{pmatrix} \mathbf{s} & -1 \end{pmatrix}}_t \left(\begin{pmatrix} \mathbf{B} \\ \mathbf{sB} + \mathbf{e} \end{pmatrix} + x \mathbf{I} \right) \approx xt$$

fully **homomorphic** encryption

$$\begin{array}{c} \boxed{t} \\ \text{sk} \end{array} \quad \boxed{\mathbf{A}} \quad \approx \quad x \quad \boxed{t}$$

$\mathbf{enc}(\text{sk}, x)$

- 2 \mathbf{A} hides x
- 3 \mathbf{A} small

$$\underbrace{(\mathbf{s} \quad -1)}_t \mathbf{G} \cdot \mathbf{G}^{-1} \left(\left(\begin{array}{c} \mathbf{B} \\ \mathbf{sB} + \mathbf{e} \end{array} \right) + x \mathbf{I} \right) \approx xt$$

fully **homomorphic** encryption

$$\underbrace{\boxed{t}}_{\text{sk}} \quad \boxed{\mathbf{A}} \approx x \quad \boxed{t}$$

$\mathbf{enc}(\text{sk}, x)$

- 2 \mathbf{A} hides x
- 3 \mathbf{A} small

$$\underbrace{\overbrace{(\mathbf{s} \quad -1)}^{\text{new } t}}_t \mathbf{G} \cdot \mathbf{G}^{-1} \left(\left(\begin{pmatrix} \mathbf{B} \\ \mathbf{sB} + \mathbf{e} \end{pmatrix} + x\mathbf{I} \right) \right) \approx xt$$

fully **homomorphic** encryption

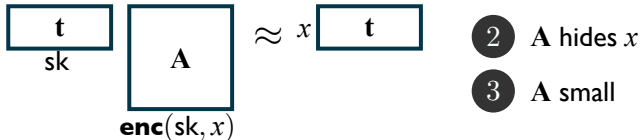
$$\underbrace{\boxed{t}}_{\text{sk}} \quad \boxed{\mathbf{A}} \approx x \quad \boxed{t}$$

$\mathbf{enc}(\text{sk}, x)$

- 2 \mathbf{A} hides x
- 3 \mathbf{A} small

$$\underbrace{\overbrace{(\mathbf{s} \quad -1)}^{\text{new } t}}_t \mathbf{G} \cdot \mathbf{G}^{-1} \left(\left(\begin{pmatrix} \mathbf{B} \\ \mathbf{sB} + \mathbf{e} \end{pmatrix} + x\mathbf{G} \right) \right) \approx xt\mathbf{G}$$

fully **homomorphic** encryption



$$\underbrace{(s \quad -1) \mathbf{G}}_t \cdot \mathbf{G}^{-1} \left(\left(\begin{array}{c} \mathbf{B} \\ \mathbf{sB} + \mathbf{e} \end{array} \right) + x\mathbf{G} \right) \approx xt$$

small, small, ...

small, small, ...

$$- \mathbf{G}^{-1}(\mathbf{M}_1)\mathbf{G}^{-1}(\mathbf{M}_2) \Rightarrow \text{small} \approx \text{small}^{\deg(f)}$$

small, small, ...

- $\mathbf{G}^{-1}(\mathbf{M}_1)\mathbf{G}^{-1}(\mathbf{M}_2) \Rightarrow \text{small} \approx \text{small}^{\deg(f)}$
- $\mathbf{G}^{-1}(\mathbf{M}_1\mathbf{G}^{-1}(\mathbf{M}_2)) \Rightarrow \text{small} \approx \text{small}^{\log \deg(f)}$

small, small, ...

circuit



intermediate \times intermediate

$$\text{small}_{\text{output}} = n^{\text{depth}} \cdot \text{small}_{\text{input}}$$

small, small, ...

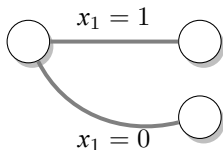
circuit



intermediate \times intermediate

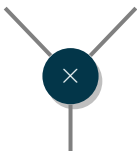
$$\text{small}_{\text{output}} = n^{\text{depth}} \cdot \text{small}_{\text{input}}$$

branching program



small, small, ...

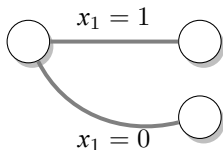
circuit



intermediate \times intermediate

$$\text{small}_{\text{output}} = n^{\text{depth}} \cdot \text{small}_{\text{input}}$$

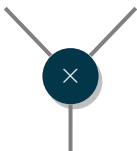
branching program



intermediate \times input

small, small, ...

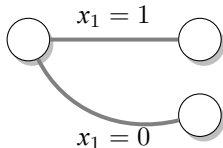
circuit



intermediate \times intermediate

$$\text{small}_{\text{output}} = n^{\text{depth}} \cdot \text{small}_{\text{input}}$$

branching program

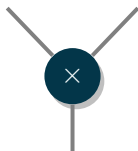


intermediate \times input

$$\text{small}_{\text{output}} = n \cdot \text{length} \cdot \text{small}_{\text{input}}$$

small, small, ...

circuit
depth $O(\log n)$

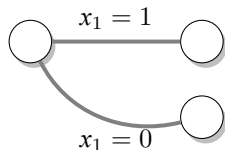


intermediate \times intermediate

$$\text{small}_{\text{output}} = n^{\text{depth}} \cdot \text{small}_{\text{input}}$$



branching program
length $\text{poly}(n)$

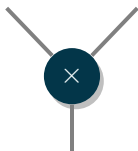


intermediate \times input

$$\text{small}_{\text{output}} = n \cdot \text{length} \cdot \text{small}_{\text{input}}$$

small, small, ...

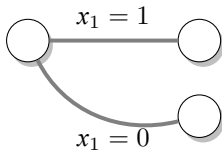
circuit
depth $O(\log n)$



$n^{O(\log n)}$ blow-up

\subseteq

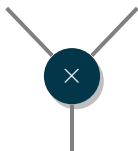
branching program
length $\text{poly}(n)$



$\text{poly}(n)$ blow-up

small, small, ...

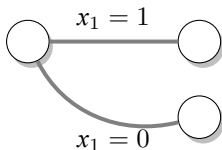
circuit
depth $O(\log n)$



$n^{O(\log n)}$ blow-up

\subseteq

branching program
length $\text{poly}(n)$



$\text{poly}(n)$ blow-up

log-depth circuits with **polynomial** hardness [BV14, AP14, GVW13]

eigenvectors, revisited

lemma I. $\mathbf{t} \cdot \mathbf{A}_i = x_i \mathbf{t} \Rightarrow \mathbf{t} \cdot \underbrace{\mathbf{A}_f}_{f(\mathbf{A}_1, \dots, \mathbf{A}_n)} = f(x) \mathbf{t}$

eigenvectors, revisited

lemma 1. $\mathbf{t} \cdot (\mathbf{A}_i - x_i \mathbf{I}) = \mathbf{0} \Rightarrow \mathbf{t} \cdot (\mathbf{A}_f - f(x) \mathbf{I}) = \mathbf{0}$

for any polynomial f , $x = (x_1, \dots, x_n)$

eigenvectors, revisited

lemma I. $\mathbf{t} \cdot (\mathbf{A}_i - x_i \mathbf{I}) = \mathbf{0} \Rightarrow \mathbf{t} \cdot (\mathbf{A}_f - f(x) \mathbf{I}) = \mathbf{0}$

lemma II. $\forall \mathbf{A}_i$

$$[\mathbf{A}_1 - x_1 \mathbf{I} \mid \cdots \mid \mathbf{A}_n - x_n \mathbf{I}] \quad \mathbf{A}_f - f(x) \mathbf{I}$$

[GSW13, BGG+14, GVW15, BCTW16, MPI2]

eigenvectors, revisited

lemma I. $\mathbf{t} \cdot (\mathbf{A}_i - x_i \mathbf{I}) = \mathbf{0} \Rightarrow \mathbf{t} \cdot (\mathbf{A}_f - f(x) \mathbf{I}) = \mathbf{0}$

lemma II. $\forall \mathbf{A}_i, \forall x, \forall f, \exists \mathbf{H}_{f,x}$

$$[\mathbf{A}_1 - x_1 \mathbf{I} \mid \cdots \mid \mathbf{A}_n - x_n \mathbf{I}] \cdot \mathbf{H}_{f,x} = \mathbf{A}_f - f(x) \mathbf{I}$$

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[GSW13, BGG+14, GVW15, BCTW16, MPI2]

claim. lemma II \Rightarrow lemma I

proof. multiply both sides by \mathbf{t}

eigenvectors, revisited

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proof. handle $+$ and \times

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proof. handle + and \times

$$[\mathbf{A}_1 - x_1 \mathbf{I} \mid \mathbf{A}_2 - x_2 \mathbf{I}] \underbrace{\left(\begin{array}{c} \\ \\ \end{array} \right)}_{\mathbf{H}_{+,x_1,x_2}} = (\mathbf{A}_1 + \mathbf{A}_2) - (x_1 + x_2) \mathbf{I}$$

eigenvectors, revisited

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proof. handle $+$ and \times

$$[\mathbf{A}_1 - x_1 \mathbf{I} \mid \mathbf{A}_2 - x_2 \mathbf{I}] \underbrace{\begin{pmatrix} \mathbf{A}_2 \end{pmatrix}}_{\mathbf{H}_{\times, x_1, x_2}} = \mathbf{A}_1 \mathbf{A}_2 - x_1 x_2 \mathbf{I}$$

eigenvectors, revisited

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eigenvectors, revisited*

lemma I. $\mathbf{t} \cdot \mathbf{A}_i = x_i \mathbf{t} \Rightarrow \mathbf{t} \cdot \underbrace{\mathbf{A}_f}_{f(\mathbf{A}_1, \dots, \mathbf{A}_n)} = f(x) \mathbf{t}$

lemma II. $\forall \mathbf{A}_i, \forall x, \forall f, \exists \mathbf{H}_{f,x}$

$$[\mathbf{A}_1 - x_1 \mathbf{I} \mid \cdots \mid \mathbf{A}_n - x_n \mathbf{I}] \cdot \mathbf{H}_{f,x} = \mathbf{A}_f - f(x) \mathbf{I}$$

$$\boxed{\mathbf{A}_i}, \boxed{\mathbf{I}} \mapsto \boxed{\mathbf{A}_i}, \boxed{\mathbf{G}}$$

eigenvectors, revisited*

lemma I. $\mathbf{t} \cdot \mathbf{A}_i = x_i \mathbf{t} \Rightarrow \mathbf{t} \cdot \underbrace{\mathbf{A}_f}_{f(\mathbf{A}_1, \dots, \mathbf{A}_n)} = f(x) \mathbf{t}$

lemma II*. $\forall \mathbf{A}_i, \forall x, \forall f, \exists$ small $\mathbf{H}_{f,x}$

$$[\mathbf{A}_1 - x_1 \mathbf{G} \mid \cdots \mid \mathbf{A}_n - x_n \mathbf{G}] \cdot \mathbf{H}_{f,x} = \mathbf{A}_f - f(x) \mathbf{G}$$

$$\boxed{\mathbf{A}_i}, \boxed{\mathbf{I}} \mapsto \boxed{\mathbf{A}_i}, \boxed{\mathbf{G}}$$

eigenvectors, revisited*

lemma I*. $\mathbf{t} \cdot \mathbf{A}_i \approx x_i \mathbf{t} \cdot \mathbf{G} \Rightarrow \mathbf{t} \cdot \underbrace{\mathbf{A}_f}_{f(\mathbf{A}_1, \dots, \mathbf{A}_n)} \approx f(x) \mathbf{t} \cdot \mathbf{G}$

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$$[\mathbf{A}_1 - x_1 \mathbf{G} \mid \dots \mid \mathbf{A}_n - x_n \mathbf{G}] \cdot \mathbf{H}_{f,x} = \mathbf{A}_f - f(x) \mathbf{G}$$

corollary. small $\mathbf{H}_{f,x} \Rightarrow$ robust to noise

eigenvectors, revisited*

$$\text{lemma I}^*. \mathbf{t} \cdot \mathbf{A}_i \approx x_i \mathbf{t} \cdot \mathbf{G} \Rightarrow \mathbf{t} \cdot \underbrace{\mathbf{A}_f}_{f(\mathbf{A}_1, \dots, \mathbf{A}_n)} \approx f(x) \mathbf{t} \cdot \mathbf{G}$$

lemma II*. $\forall \mathbf{A}_i, \forall x, \forall f, \exists$ small $\mathbf{H}_{f,x}$

$$[\mathbf{A}_1 - x_1 \mathbf{G} \mid \dots \mid \mathbf{A}_n - x_n \mathbf{G}] \cdot \mathbf{H}_{f,x} = \mathbf{A}_f - f(x) \mathbf{G}$$

corollary. small $\mathbf{H}_{f,x} \Rightarrow$ robust to noise

$$(\mathbf{s}[\mathbf{A}_1 - x_1 \mathbf{G} \mid \dots \mid \mathbf{A}_n - x_n \mathbf{G}] + \mathbf{e}) \cdot \mathbf{H}_{f,x} \approx \mathbf{s}(\mathbf{A}_f - f(x) \mathbf{G})$$

eigenvectors, revisited*

$$\text{lemma I}^*. \quad t \cdot A_i \approx x_i t \cdot G \Rightarrow t \cdot \underbrace{A_f}_{f(A_1, \dots, A_n)} \approx f(x) t \cdot G$$

lemma II*. $\forall A_i, \forall x, \forall f, \exists$ small $H_{f,x}$

$$[A_1 - x_1 G \mid \dots \mid A_n - x_n G] \cdot H_{f,x} = A_f - f(x)G$$

proof. handle + and \times

eigenvectors, revisited*

lemma I*. $t \cdot \mathbf{A}_i \approx x_i t \cdot \mathbf{G} \Rightarrow t \cdot \underbrace{\mathbf{A}_f}_{f(\mathbf{A}_1, \dots, \mathbf{A}_n)} \approx f(x) t \cdot \mathbf{G}$

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$$[\mathbf{A}_1 - x_1 \mathbf{G} \mid \mathbf{A}_2 - x_2 \mathbf{G}] \underbrace{\begin{pmatrix} \mathbf{I} \\ \mathbf{I} \end{pmatrix}}_{\text{small}} = (\mathbf{A}_1 + \mathbf{A}_2) - (x_1 + x_2) \mathbf{G}$$

eigenvectors, revisited*

lemma I*. $t \cdot \mathbf{A}_i \approx x_i t \cdot \mathbf{G} \Rightarrow t \cdot \underbrace{\mathbf{A}_f}_{f(\mathbf{A}_1, \dots, \mathbf{A}_n)} \approx f(x) t \cdot \mathbf{G}$

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$$[\mathbf{A}_1 - x_1 \mathbf{G} \mid \mathbf{A}_2 - x_2 \mathbf{G}] \underbrace{\begin{pmatrix} \mathbf{A}_2 \\ x_1 \mathbf{I} \end{pmatrix}}_{\text{small?}} = \mathbf{A}_1 \mathbf{A}_2 - x_1 x_2 \mathbf{G}$$

eigenvectors, revisited*

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$$[A_1 - x_1 G \mid \dots \mid A_n - x_n G] \cdot H_{f,x} = A_f - f(x)G$$

proof. handle $+$ and \times

$$[A_1 - x_1 G \mid A_2 - x_2 G] \underbrace{\begin{pmatrix} G^{-1}(A_2) \\ x_1 I \end{pmatrix}}_{\text{small}} = A_1 G^{-1}(A_2) - x_1 x_2 G$$

eigenvectors, revisited*

lemma I*. $t \cdot \mathbf{A}_i \approx x_i t \cdot \mathbf{G} \Rightarrow t \cdot \underbrace{\mathbf{A}_f}_{f(\mathbf{A}_1, \dots, \mathbf{A}_n)} \approx f(x) t \cdot \mathbf{G}$

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$$[\mathbf{A}_1 - x_1 \mathbf{G} \mid \mathbf{A}_2 - x_2 \mathbf{G}] \underbrace{\begin{pmatrix} \mathbf{G}^{-1}(\mathbf{A}_2) \\ x_1 \mathbf{I} \end{pmatrix}}_{\text{small}} = \underbrace{\mathbf{A}_1 \mathbf{G}^{-1}(\mathbf{A}_2)}_{\mathbf{A}_x} - x_1 x_2 \mathbf{G}$$

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applications.

fully homomorphic enc [GSW]

attribute-based enc [BGGHNSVV]

fully homomorphic sig [GVW]

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applications.

	A_f	$H_{f,x}$
fully homomorphic enc [GSW]	eval output	correctness

attribute-based enc [BGGHNSVV]

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applications.

	A_f	$H_{f,x}$
fully homomorphic enc [GSW]	eval output	correctness
attribute-based enc [BGGHNSVV]	keygen	decryption
fully homomorphic sig [GVW]	verification	homomorphic sign

conclusion

today. lattices \Rightarrow **encrypted computation**

$$[\mathbf{A}_1 - x_1 \mathbf{G} \mid \cdots \mid \mathbf{A}_n - x_n \mathbf{G}] \cdot \mathbf{H}_{f,x} = \mathbf{A}_f - f(x) \mathbf{G}$$

conclusion

today. lattices \Rightarrow **encrypted computation**

$$[\mathbf{A}_1 + x_1 \mathbf{G} \mid \cdots \mid \mathbf{A}_n + x_n \mathbf{G}] \cdot \mathbf{H}_{f,x} = \mathbf{A}_f + f(x) \mathbf{G}$$

conclusion

today. lattices \Rightarrow **encrypted computation**

$$[\mathbf{A}_1 - x_1 \mathbf{G} \mid \cdots \mid \mathbf{A}_n - x_n \mathbf{G}] \cdot \mathbf{H}_{f,x} = \mathbf{A}_f - f(x) \mathbf{G}$$

today. lattices \Rightarrow **FHE** for circuits with **dec** \approx $\langle \text{sk}, \text{ct} \rangle$

conclusion

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today. lattices \Rightarrow **FHE** for circuits with **dec** \approx $\langle \text{sk}, \text{ct} \rangle$

“ XXX for fhe.dec \Rightarrow XXX for circuits ” [GVW12,GKPVZ13,GVW15]

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starting point for **obfuscation – tomorrow**

conclusion

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$$[\mathbf{A}_1 - x_1 \mathbf{G} \mid \cdots \mid \mathbf{A}_n - x_n \mathbf{G}] \cdot \mathbf{H}_{f,x} = \mathbf{A}_f - f(x) \mathbf{G}$$

tue. obfuscation

wed. another way to encode computation into lattices

[GGH15, KW16, CC17, GKW17, WZ17, GKW18, CVW18, ...]

conclusion

today. lattices \Rightarrow **encrypted computation**

$$[\mathbf{A}_1 - x_1 \mathbf{G} \mid \cdots \mid \mathbf{A}_n - x_n \mathbf{G}] \cdot \mathbf{H}_{f,x} = \mathbf{A}_f - f(x) \mathbf{G}$$

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wed. another way to encode computation into lattices

thu. MPC, LWE, FHE

conclusion

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$$[\mathbf{A}_1 - x_1 \mathbf{G} \mid \cdots \mid \mathbf{A}_n - x_n \mathbf{G}] \cdot \mathbf{H}_{f,x} = \mathbf{A}_f - f(x) \mathbf{G}$$

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thu. MPC, LWE, FHE

fri. quantum crypto

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tue. obfuscation

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// thank you & enjoy!