

Optimal Privacy-Constrained Mechanisms

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Motivation

Introduce and analyze a Bayesian measure of privacy loss.

- Most work on differential privacy (Dwork et al. '06) is “prior-free”
- From an **outsider**'s perspective, the realized outcome of a DP mechanism does not reveal much about any individual participant's type

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- But to implement this, types have to be reported
- We might worry about the **designer** knowing too much
- Our approach: mechanism design under a privacy constraint that *limits how much information the principal can collect from the agents*

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- 6 **Principal constrained by $I(\mathbb{M}) \leq \kappa$ with κ exogenously given**

Discussion

- Definition equivalent to **MI between types and messages**:
 - ▶ early version of Xiao ('13) considers MI as a cost to each agent
 - ▶ we take the paternalistic viewpoint of a **regulator** but do not directly model agent preferences for privacy
 - ▶ alternatively, each agent participates only if constraint is met
- Above measure of privacy loss takes average across different messages:
 - ▶ more stringent “ex-post” notion requires $D(F(\cdot | m) || F) \leq \kappa, \forall m$
 - ▶ results similar; focus on ex-ante case here
- Related issue of how to aggregate privacy loss across multiple agents:
 - ▶ paper studies an application with only one agent

Screening Environment

Focus on the monopolistic screening model of Mussa-Rosen ('78).

- A seller sells some quantity/quality $q \geq 0$ to a buyer for payment p
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- In our model, seller maximizes profit subject to privacy. That is,

$$\max \mathbb{E}_m [p(m) - c(q(m))] \quad s.t. \quad \mathbb{E}_m [D(F(\cdot | m) || F)] \leq \kappa$$

Coarse Revelation

Main Result

Given $0 < \kappa < \infty$. There exists an optimal privacy-constrained mechanism \mathbb{M} , where the set of types $[\underline{\theta}, \bar{\theta}]$ is partitioned into finitely many intervals, and in equilibrium each type truthfully reports its interval.

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Further properties:

- privacy constraint binds in any optimal mechanism
- if κ small, exactly two intervals used

Why Intervals?

Several papers (e.g. Bergemann et al.) derived optimality of intervals by assuming upper bound on number of messages. For us,

- 1 First remove “redundant” messages: If two messages lead to same outcome, combine them into a single message
 \implies posterior belief is *averaged*, implying smaller privacy loss
- 2 Types that send different messages partition the type space
- 3 By single-crossing property, each partition is convex
- 4 Thus intervals – this part does not rely on specific form of KL; also extends to multiple agents with one-dimensional types

Finiteness

Where we use KL is to show finite intervals suffice.

- Technical difficulty as space of partitions is *not compact*
- We restore compactness by showing **at most one short interval**
- Otherwise, *merge two short intervals* and use saved privacy to *divide a long interval*. Profit would increase
- Intuition: “log” term in KL punishes heavily against getting precise information about even a small set of types

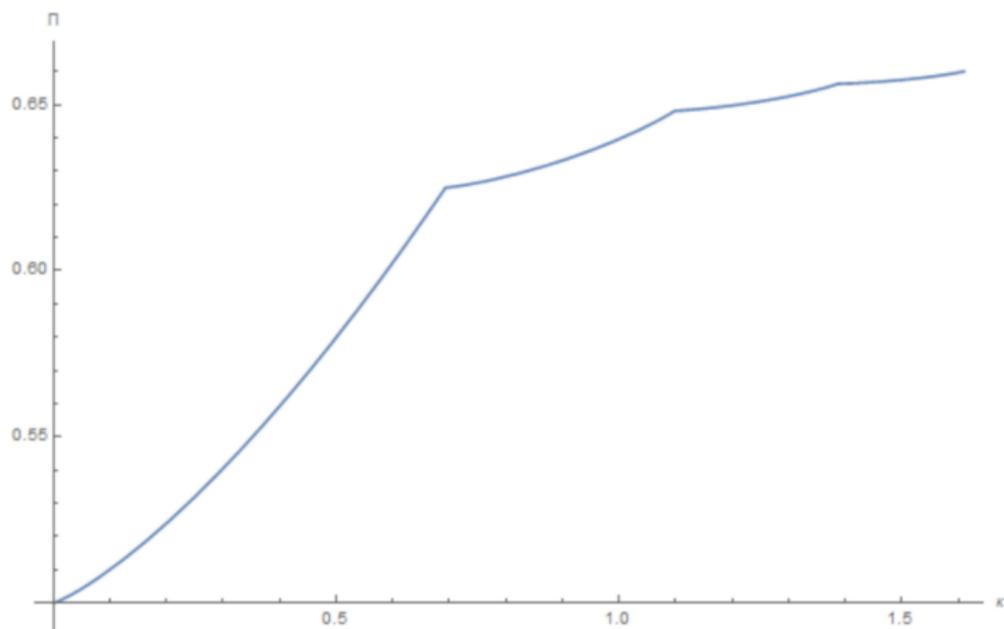
Uniform Case

Consider special case with uniform types. Can show “ordering” of intervals do not matter for profit and privacy measure.

Characterization

With uniform prior, for any κ , the optimal privacy-constrained mechanism partitions $[\underline{\theta}, \bar{\theta}]$ into $n - 1$ *equally long* intervals and 1 *shorter* interval, such that the privacy constraint is exhausted.

Profit Frontier



plotted for $\theta \sim U [1, 2]$

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Comparative Statics w.r.t. κ

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- 1 Profit from a κ -constrained optimal mechanism increases in κ
- 2 Buyer surplus is maximized (resp. minimized) with full (resp. no) privacy
- 3 If prior density $f(\theta)$ decreases, no privacy maximizes total welfare

Recap

- Bayesian privacy measure: how much principal learns via mechanism
⇒ Coarse menu offered in the form of interval partition
- Implementation: where does the prior come from?
- Multiple agents: how to aggregate privacy?
- Dynamic mechanisms?

Thank You!