

# VerITAS: Verifying Image Transformations at Scale

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# These look like prizewinning photos. They're AI fakes.

Artificially generated images of real-world news events proliferate on stock image sites, blurring truth and fiction

By [Will Oremus](#) and [Pranshu Verma](#)

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AI-GENERATED FAKE PHOTO



AI-GENERATED FAKE PHOTO



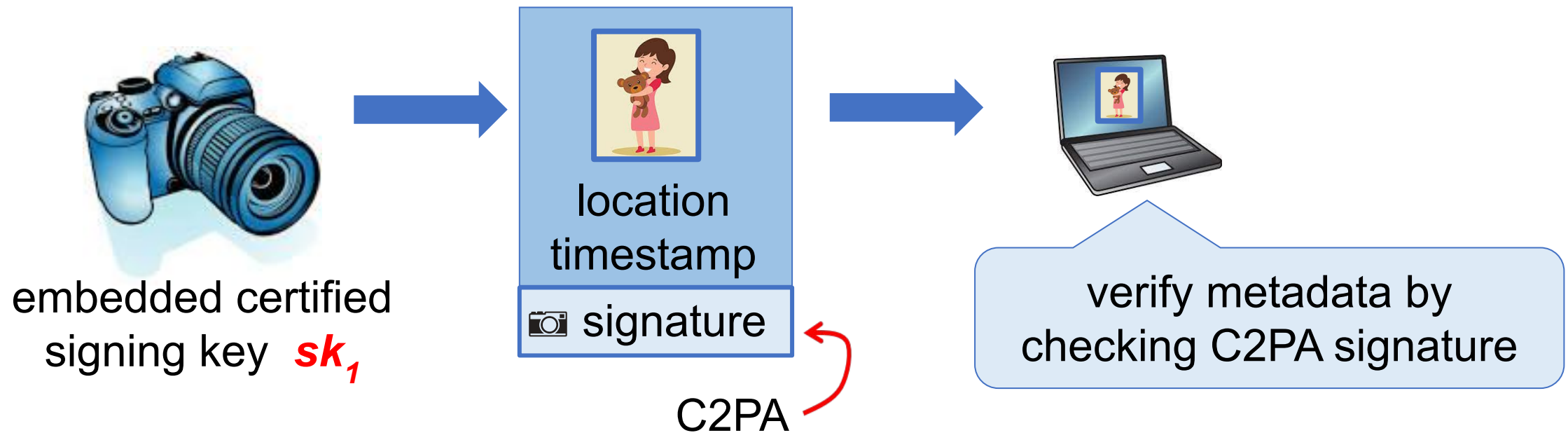
AI-GENERATED FAKE PHOTO



# C2PA: A Content Provenance Standard

Nikon, Canon, Sony eye tamper-resistant digital signatures to combat deepfakes

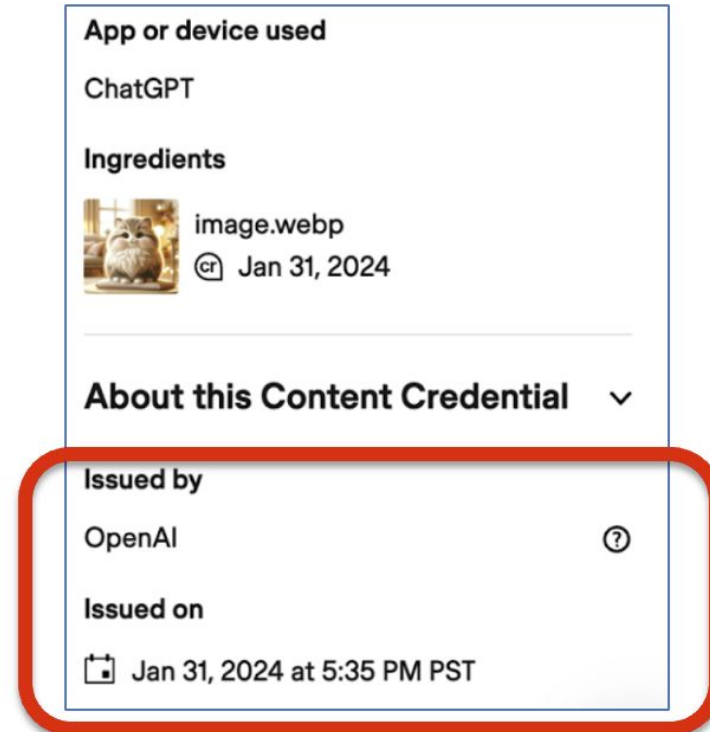
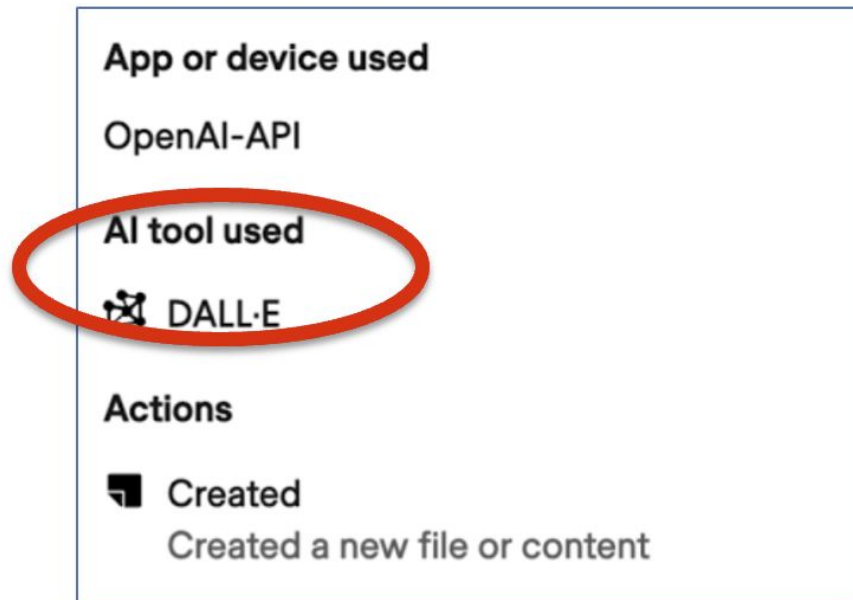
Partnership for greater trust in digital photography: Leica and Content Authenticity Initiative



# Not just news organizations...

## C2PA in DALL·E 3

C2PA standard, OpenAI's implementation, and C2PA metadata



# Is this a cat-and-mouse game?

The picture-of-picture attack:

Many other challenges

- (1) Key extraction and revocation (PKI)
- (2) Privacy → group signatures
- (3) GPS spoofing

Now every v

- Is this
- Can attacker defeat the filter?

[see Chimera, Riva, et al., [Harmful Use of the C2PA](#), 2022 ]



C2PA camera



C2PA Image

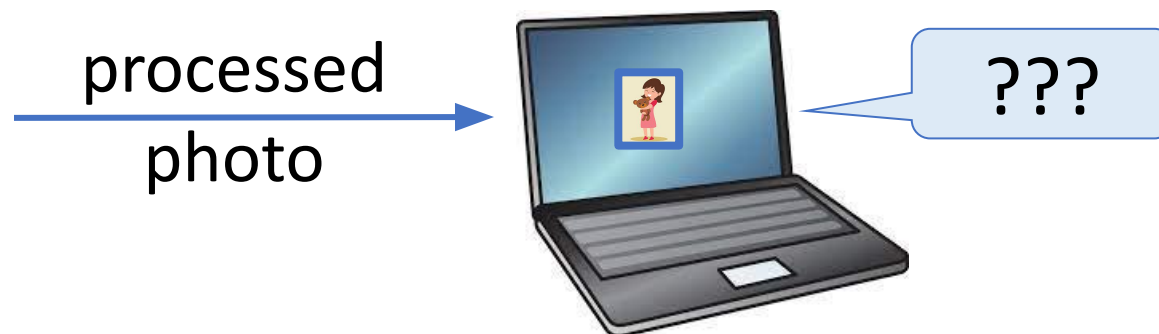


# A Problem: Post-Processing

Newspapers often process photos before publication

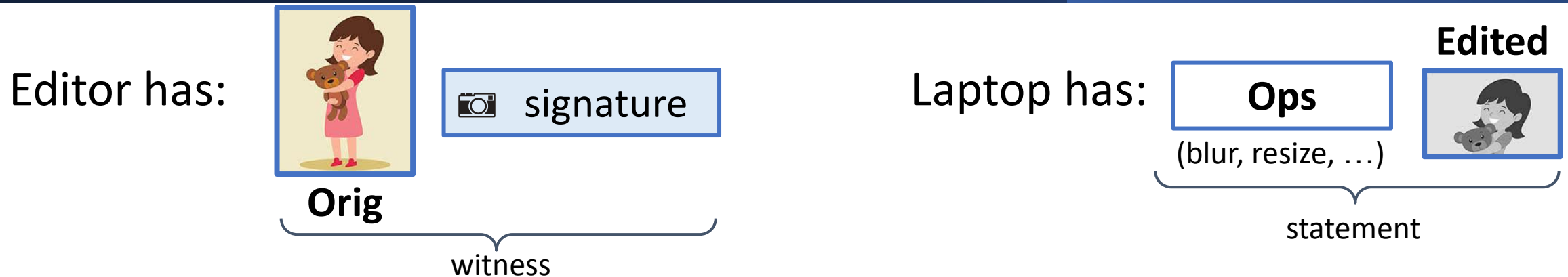
- At minimum, need to resize (90 MB → 8 MB)
- Allowable operations from the *Associated Press*:  
cropping, grayscale, blurring...

**Problem:** browser cannot verify the C2PA signature of a processed photo



The solution proposed by C2PA is not ideal ... can we do better?

# A Cryptographic Solution: zk-SNARKs!



Editor creates a proof  $\pi$  that:

I know **(*Orig*, *sig*)** such that:

1. ***sig*** is a valid signature on ***Orig***
2. ***Edited*** is the result of applying ***Ops*** to ***Orig***
3.  $\text{metadata}(\mathbf{Edited}) = \text{metadata}(\mathbf{Orig})$

Do we need  
ZK ZK?

$\Rightarrow$  Laptop verifies  $\pi$  and shows metadata to user

# How to prove?

$\pi$

I know **(*Orig*, *sig*)** such that:

1. ***sig*** is a valid signature on ***Orig***
2. ***Edited*** is the result of applying ***Ops*** to ***Orig***
3.  $\text{metadata}(\mathbf{Edited}) = \text{metadata}(\mathbf{Orig})$

?



# Verifying Edits in a SNARK Prover

- PhotoProof (Naveh and Tromer, 2016): a few minutes to generate photo editing proofs for 128 x 128 pixel image
- New tools enable faster development and bigger statements!
  - Plonky2 library (“Plonk PIOP” + FRI PCS)
    - Write arithmetic circuit  $C_{\text{edit}}$  s.t.  $C_{\text{edit}}(\mathbf{Orig}) = \mathbf{Edited}$

Our work: proof for a 6000 x 4000 image using Plonky2

- resize, crop, grayscale, blur  $\rightarrow$  proof gen. time  $\leq 4$  minutes
- Proof size:  $\approx 100$  KB ( $\ll$  image size), verification time: 0.7 sec  
(can shrink proof with recursion)

# Verifying Signatures in a SNARK Prover

$\pi$

I know (***Orig***, ***sig***) such that:

1. ***sig*** is a valid signature on ***Orig***
2. ***Edited*** is the result of applying ***Ops*** to ***Orig***
3.  $\text{metadata}(\textbf{Edited}) = \text{metadata}(\textbf{Orig})$



**Problem:** the SNARK proof must check that a pair (***Orig***, ***h***) satisfies  $\mathbf{h} = \text{hash}(\textbf{Orig})$ ...but ***Orig*** is 90MB!

# Verifying Signatures in a SNARK Prover

## Attempt 1

### Verify Hash

I know (*Orig*, *hash*) such that:

$$\text{hash} = \text{SHA256}(\text{Orig})$$

Too slow for  
90 MB!

## Attempt 2

### Verify Hash

I know (*Orig*, *hash*) such that:

$$\text{hash} = \text{Poseidon}(\text{Orig})$$

SNARK-friendly hash  
...but still too slow for  
90MB!

We propose two methods:

### (1) Verify Lattice Hash

I know (*Orig*, *hash*) such that:  
*hash* =  
 $\text{Poseidon}(\text{LatticeHash}(\text{Orig}))$

good for camera and prover

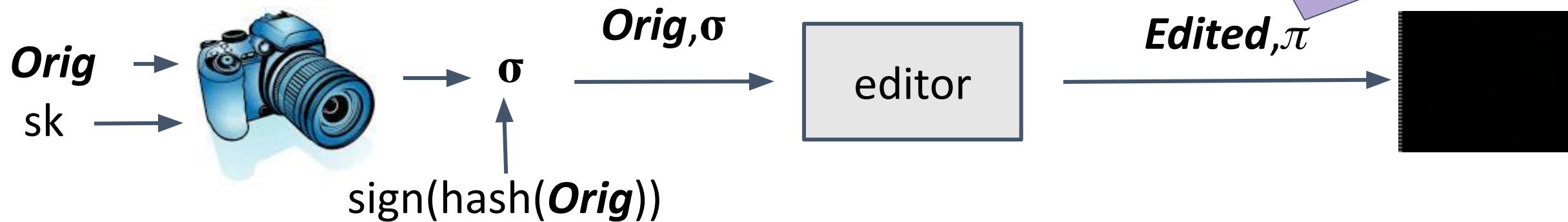
### (2) Verify a PCS

I know (*Orig*, *hash*) such that:  
*hash* =  $\text{PCS}(\text{Orig})$

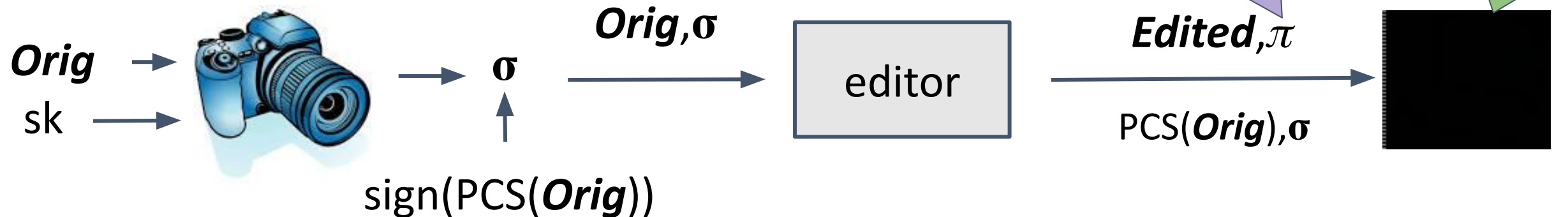
great for prover!

# Verifying Signatures in a SNARK Prover

## Method 1: lattice hash

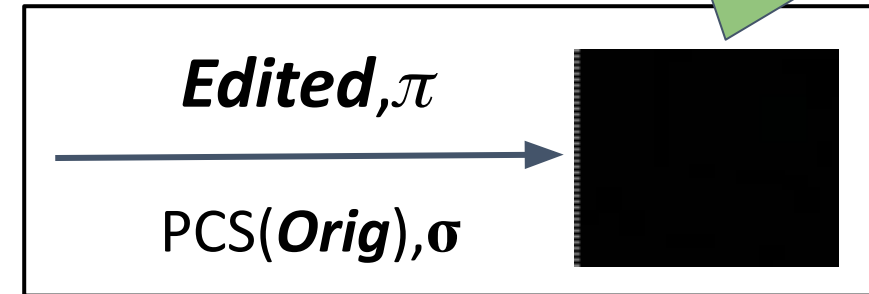


## Method 2: use polynomial commitment scheme (PCS)



# How is Method 2 secure?

PLONK proof  $\pi$  proves that  $C_{\text{edit}}(\mathbf{Orig}') = \mathbf{Edited}$   
where  $\mathbf{Orig}'$  is provided as witness data



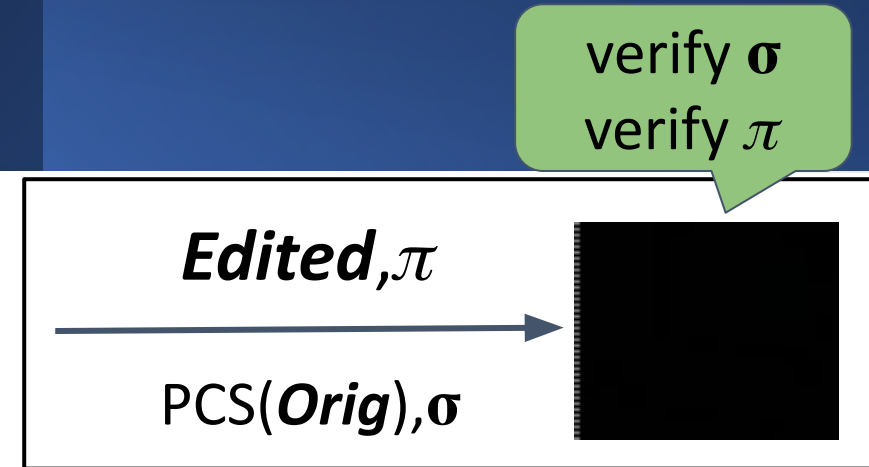
Problem: what if  $\mathbf{Orig} \neq \mathbf{Orig}'$  ??

→ Then edited image is unrelated to camera's image

Solution: non-black-box use of PLONK!

# How is Method 2 secure?

PLONK proof  $\pi$  proves that  $C_{\text{edit}}(\mathbf{Orig}') = \mathbf{Edited}$   
where  $\mathbf{Orig}'$  is provided as witness data



Partial explanation of how to produce PLONK proof  $\pi$ :

1. Encode  $C_{\text{edit}}$  execution tableaux as a polynomial

$T(x)$



Some  $T(x)$  evaluations  
encode the witness  $\mathbf{Orig}'$

2. Generate proof:

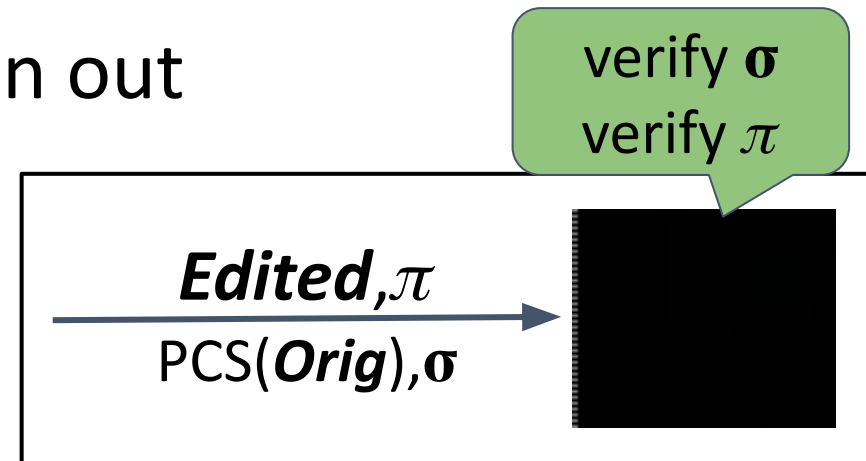
- Compute  $\mathbf{com}_T = \text{PCS}(T(x))$
- Prove gates evals in tableaux are correct
- Prove circuit wiring in tableaux is correct

This works even if tableaux is  
committed via multiple polynomials!



# Tradeoffs of Signing PCS vs. Lattice Hash

Signing PCS(***Orig***) takes signature verification out of the SNARK circuit



... but computing a PCS commitment is  
not feasible on a commercial camera

- Suitable for a cloud AI image generator
- Can be offloaded to an untrusted server

# Verifying Signatures in a SNARK Prover

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# Conclusions

Succinct proofs have become practical and easy to use

- An amazing success of theory of CS
- Development driven by blockchain  
but many non-blockchain applications

C2PA: a playground for many cryptographic techniques

- Many challenges to explore...