

#### **QUANTUM PROGRAM VERIFICATION** AN AUTOMATA-BASED APPROACH

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Games and Equilibria in System Design and Analysis

#### Why Quantum Computing Is Important?



#### **Promises:**

- Solve conventional unsolvable problems.
- Example: break cryptography
- ▶ Algorithms for solving practical problems are under fast developing:
	- Machine learning
	- Optimization
	- Quantum chemistry





The quantum computer Jiuzhang manipulates light via a complex arrangement of optical devices (shown HANSEN ZHONG

# Quantum Software Stack



- Classical software handles tasks like control flow, algorithm design, and data preprocessing.
- The synergy between quantum and classical systems is essential for potential applications.

 $U1^{\theta}$ ,  $U2^{\theta_1,\theta_2}$ ,  $U3^{\theta_1,\theta_2,\theta_3}$ , CX Ion **IBM** 



 $R_x^{\theta}$ ,  $R_y^{\theta}$ ,  $R_z^{\theta}$ ,

3

# Quantum Software Stack



- Quantum computers are not standalone; they require classical software support.
- Classical software handles tasks like control flow, algorithm design, and data preprocessing.
- The synergy between quantum and classical systems is essential for potential applications.



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#### Quantum Software Correctness



5

- **Challenges in Software Development**
	- Software complexity grows, making correctness harder to ensure.
	- Debugging costs exceed half of software development expenses.
- **Quantum Software Development**
	- Traditional methods struggle due to quantum's probabilistic nature.
	- Quantum states collapse upon observation, hampering traditional testing.

#### 4**Formal Verification**

• Provides a highly effective means of ensuring the quality of quantum software.

# Verification via Examples



 $\triangleright$  Pre: The default initial state.



 $\triangleright$  Post: Found the hidden string (110 in this case).

#### **Bernstein Vazirani Algorithm**

# Verification via Examples



4Post: |110−⟩.

#### **Bernstein Vazirani Algorithm**

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# Screenshots of AutoQ





# Screenshots of AutoQ





# **Outline**



- **Quantum Background**
- Quantum Circuit Verification
- Quantum Program Verification

### A 3-Bit Classical State





# A 3-Qubit Quantum State







## Tree as a Quantum State



#### ▶ A 3-bit quantum state



#### Quantum Gate and Tree Transformation



#### An example of apply  $X$  gate (negation) on qubit  $x_1$ .



#### Quantum Gate and Tree Transformation



An example of apply H gate on qubit  $x_1$ .



#### Quantum Gate and Tree Transformation



An example of apply CX gate on control qubit  $x_1$  and target qubit  $x_2$ .



# Quantum Parallelism



#### • One gate updates an exponential number of classical states





18



# Quantum Circuit





# Quantum Simulation





Multi-Terminal Binary Decision Diagram (MTBDD)



The EPR circuit



 $x_1$ 

0

1



# Quantum Simulation





Multi-Terminal Binary Decision Diagram (MTBDD)

In classical verification

BDDs encodes a set of states In quantum, it encodes one state - how to encode a set of state?



The EPR circuit

Compress







# **Outline**



- ▶ Quantum Background
- 4**Quantum Circuit Verification**
- 4Quantum Program Verification

## Classical Hoare triple



 $\triangleright$  For any predicates P and Q and any program S,

{P} S {Q} Precondition  $\{P\}$   $\{ \}$   $\{ \}$  Postcondition

says that if  $S$  is started in (a state satisfying) P, then it terminates in Q.

# Quantum Circuit Verification

#### Need a symbolic representation of a set of quantum states (trees).

# $\therefore$   $\triangle$   $\{P\}$  C  $\{Q\}$   $\triangle$

Set of trees  $\rightarrow$  Regular tree language (Tree automata) From automata theory: Set of words  $\rightarrow$  Regular language (Finite automata)























#### Three Components for Symbolic Verification

- **1. Symbolic representation of sets of states**
- 2. Algorithm to compute the post image
- 3. Algorithm to check containment

#### **Examples: Tree Automata** Encoding of Quantum States

 $\chi_1$ 

 $\chi_2$ 



 $\triangleright$  This TA accepts all 3-qubit basis quantum states {|000⟩, |001⟩, |010⟩, |011⟩,

|100⟩, |101⟩, |110⟩, |111⟩}



BDD + non-deterministic branching

#### TA as Compact Representation of Quantum States **ACADEMIA SINICA**



- $\triangleright$  This TA accepts all  $2<sup>n</sup>$  basis states.
- $\triangleright$  # of transitions:  $3n+1$

Why it can be some compact? Merge shared structures.

## How about this



$$
\left\{\tfrac{1}{\sqrt{2}}(|0b_2b_3\ldots b_n\rangle+|1\bar{b}_2\bar{b}_3\ldots\bar{b}_n\rangle)\,\,\big|\,\,b_2b_3\ldots b_n\in\mathbb{B}^{n-1}\right\}
$$



- $\blacktriangleright$  A common pattern of reachable set of quantum states.
- $\blacktriangleright$  Need at least  $2^{n-1}$  root transitions.

#### Level Synchronized Tree Automata (under submission)



- Transitions are labeled with "choices"  $\{1\}, \{2\}, \text{ or } \{1,2\}$
- A run only allows transitions with common choice at the same level.
- Choices of transitions from one state must be disjoint.

# $\{\frac{1}{\sqrt{2}}(|00\rangle \pm |11\rangle), \frac{1}{\sqrt{2}}(|01\rangle \pm |10\rangle)\| \pm \in \{+,-\}\}\$

#### Level Synchronized Tree Automata (under submission)



- Incomparable expressiveness compared with standard TA.
- Language inclusion is decidable.

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#### Level Synchronized Tree Automata (under submission)



 $\{|0^n\rangle | n \geq 1\}$ 

Take **choice 1** to make a next level or **choice 2** to leaves

BDD + non-deterministic branching + cycle

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#### Three Components for Symbolic Verification

- 1. Symbolic representation of sets of states
- **2. Algorithm to compute the post image**
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#### **Examples of Gate Operations:** X gate on qubit 2.



 $q-(x_1)$   $\rightarrow$   $(q_0, q_0)$  $q_0-(x_2)$  +  $(q_1,q_2)$ 

 $q_1 - 1$  $q_2$  –(0)+()



 $(x_1) (q_0, q_0)$  $q_0$   $\left(x_2\right)$  +  $\left(q_2, q_1\right)$ 

 $q_1$ -



#### **Example of Gate Operations:** Z, S, T gates on qubit 1.

#### $\blacktriangleright$  Multiply the right subtree of  $x_1$  with some constant c.

$$
q - (x_1) \rightarrow (q_0^1, q_1^1) \qquad q_1^1 - (x_2) \rightarrow (q_0^2, q_1^2) \qquad q_1^2 - (x_3) \rightarrow (q_0, q_1) \qquad q_0 - (0) \rightarrow ()
$$
  
\n
$$
q - (x_1) \rightarrow (q_1^1, q_0^1) \qquad q_1^1 - (x_2) \rightarrow (q_1^2, q_0^2) \qquad q_1^2 - (x_3) \rightarrow (q_1, q_0) \qquad q_1 - (1) \rightarrow ()
$$
  
\n
$$
q_0^1 - (x_2) \rightarrow (q_0^2, q_0^2) \qquad q_0^2 - (x_3) \rightarrow (q_0, q_0)
$$
  
\n
$$
q - (x_1) \rightarrow (q_0^1, q_1^1) \qquad q_1^1 - (x_2) \rightarrow (q_0^2, q_1^2) \qquad q_1^2 - (x_3) \rightarrow (q_0, q_1) \qquad q_0 - (c \times 0) \rightarrow ()
$$
  
\n
$$
q - (x_1) \rightarrow (q_1^1, q_0^1) \qquad q_1^1 - (x_2) \rightarrow (q_0^2, q_1^2) \qquad q_1^2 - (x_3) \rightarrow (q_0, q_1) \qquad q_1 - (c \times 1) \rightarrow ()
$$
  
\n
$$
q_0^1 - (x_2) \rightarrow (q_0^2, q_0^2) \qquad q_0^2 - (x_3) \rightarrow (q_0, q_0)
$$



#### Three Components for Symbolic Verification

- 1. Symbolic representation of sets of states
- 2. Algorithm to compute the post image
- **3. Algorithm to check containment**





43

## Symbolic Extension

#### $\blacktriangleright$  Note this is different from symbolic automata



Yu-Fang Chen, Kai-Min Chung, Ondřej Lengál, Jyun-Ao Lin, Wei-Lun Tsai, AUTOQ: An Automata-Based Quantum Circuit Verifier (CAV 2023)

# **Outline**



- ▶ Quantum Background
- 4Quantum Circuit Verification
- 4**Quantum Program Verification (on-going)**

# Quantum Programs



Quantum program =

Quantum circuit + Conditional statement + Loop

 $H_1$ ;  $CX_2^1$ ; **if**  $M_1 = 0$  then  $\{X_1\}$ ; while  $M_1 = 0$  do  $\{X_1; H_1; CX_2^1\};$ 

#### Conditional Statement 央研究院 **ACADEMIA SINICA**  $H_1$ ;  $C$ )  $x_1$  $x_1$  $= 0$  then  $\{X_1\};$  $x_1$  $x_2$  $x<sub>2</sub>$  $x_2$  $\frac{-a_1}{\sqrt{2}}$   $\frac{-a_0}{\sqrt{2}}$  $rac{a_1}{\sqrt{2}}$  $rac{a_0}{\sqrt{2}}$  $rac{a_0}{\sqrt{2}}$  $rac{a_1}{\sqrt{2}}$  $\bf{0}$ 0 0  $\bf{0}$

(c)  $M_1 = 0$ 

(b) Applied  $H_1$ ;  $CX_2^1$ 

Normalize?

(d)  $M_1 = 1$ 

 $x_{2}$ 

 $\overline{0}$ 

 $x_1$ 

 $\chi_2$ 

 $a_1$ 

 $a<sub>0</sub>$ 

# Loop Statement



#### Algorithm 2: " $-X_2$ "

- 1 Pre:  $\{(a_0 | 10 \rangle + a_1 | 11 \rangle + 0 | * \rangle)\};$
- 2  $H_1$ ;  $CX_2^1$ ;
- 3 Inv:  $\left\{\frac{a_0^2}{\sqrt{2}}|00\rangle + \frac{a_1}{\sqrt{2}}|01\rangle \frac{a_1}{\sqrt{2}}|10\rangle \frac{a_0}{\sqrt{2}}|11\rangle\right\};$
- 4 while  $M_1 = 0$  do  $\{X_1; H_1; CX_2^1\};$

5 Post:  $\{(-a_1 | 10\rangle - a_0 | 11\rangle + 0 | *\rangle)\};$ 



# Reference

- ▶ An Automata-Based Framework for Verification and Bug Hunting in Quantum Circuits (PLDI 2023) https://dl.acm.org/doi/10.1145/359127
- 4 AUTOQ: An Automata-Based Quantum Circuit Verifier (CAV 2023) https://link.springer.com/chapter/10.1007/978-3-031-37709-9\_7
- ▶ Verifying Quantum Circuits with Level-Synchronized Tree Automata submission)
- 4 AutoQ 2.0: From Verification of Quantum Circuits 2 to Verification of Quantum Programs (working draft)



Summary:

• Quantum computers are not standalone; they require classical software support.

49

- Formal verification is a promising approach for ensuring quantum software quality.
- Plenty of new research problems.