# The complexity of ground states 

Zeph Landau

## Ground State condensed matter dure algorithm "-" Area Law AGSP $\begin{gathered}\text { entanglement } \\ \text { vianie } \\ \text { aet }\end{gathered}$ Local Hamiltonian

## The difficulty of understanding many-body physics



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So even describing a state requires exponential amount of information.

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- Focus on unique ground state and constant gap.

Ground states model the state of the system at low temperatures.

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Solving, classifying, and understanding the structure of the solutions of CSP's at the heart of complexity theory.

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Complexity Theory


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Complexity Theory
Condensed Matter Physics


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Complexity Theory Condensed Matter Physics



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Complexity Theory Condensed Matter Physics



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CSP constraints correspond to $H_{i}$ that are diagonal in the standard basis. In particular they all commute.

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Sure you can do it in practice . . . but can you do it in theory?

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['07, Aharonov, Gottesman, Irani, Kempe]
- Solutions to 1D systems are also hard.


## Area Law formulation

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Became known as an Area Law.
['01, Vidal, Latorre, Rico, Kitaev] Area Law formalized in terms of entanglement entropy.

- Effect on DMRG: speedup, simplification, better understanding of the heuristics used.


## Area Law in 1D systems



1D Area law proved [Hastings '07].

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['08, Cirac, Schuch, Verstraete] Example of finding a solution that satisfies the area law that is NP-hard.


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## Approximate Ground State Projection (AGSP)



Properties:

- It "approximately" projects onto one vector you want (ground state).
- It isn't too complex.

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## Current view of 1D local Hamiltonians



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## From AGSP's:

- exponential improvement on the constants for the 1D Area Law
- algorithm for 1D,
- gives insight as to what is going on,
- tools for attacking the 2D questions.


## Role of AGSP in proof of Area Law I

Two main steps:

1. Find a not very complex state that has constant overlap with the ground state.


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Both steps use AGSPs- the first is much more delicate.

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Critical threshold $D \Delta<1$.

## Role of AGSP in proving Area Law cont.

Theorem (Area Law) [Arad, Landau, Vazirani] The existence of an AGSP $K$ for which $D \Delta<1 / 2$ proves that the ground state has entropy $O(1) \log D$.

## Proof:

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AGSP definition uses notions: truncation away from cut, Chebyshev polynomials. Requires careful analysis to bound complexity. (See arxiv, find me, or future workshop).

Finding the ground state of 1D systems: solving a large convex program
finding the minimal energy state
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- A succinct description of the elements of that subspace that allows us to perform linear algebra efficiently.


## The algorithm: a bird's eye view

A sequence of spaces $S_{i}$ termed viable sets:

- all polynomial size
- all with succinct descriptions that allow efficient linear algebra,
- each containing a good approximation of the "left" side of the ground state.



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Arxiv, find me, later workshop.

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"The future ain't what it used to be." - Yogi Berra


## AGSP construction: norm reduction

Looking for low entanglement operators that look like:


Smaller $\|H\|$ would be better but we don't want to lost the local structure around the cut.
Solution: Replace $H=\sum_{i} H_{i}$ with $H^{\prime}=H_{L}+H_{1}+H_{2}+\cdots+H_{s}+H_{R}$.


## AGSP construction: Chebyshev polynomials

Chebyshev polynomials: small in an interval:


The desired AGSP is a dilation and translation of the Chebyshev polynomial:


## AGSP complexity: Entanglement rank analysis

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Total: $d^{2 \ell / s+s}$


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## AGSP complexity: Entanglement rank analysis

Problem: Too many $\left(s^{\ell}\right)$ terms in naive expansion of $\left(H^{\prime}\right)^{\ell}$.
Need to group terms in a nice way but it all works out with total entanglement increase of the same order as the single term.

## Putting things together: Area Law for $H^{\prime}$

Chebyshev $C_{\ell}\left(H^{\prime}\right)$ has $\Delta \approx e^{-O(\ell / \sqrt{s})}$ :


Entanglement analysis yields $D \approx O\left(d^{\ell / s+s}\right)$.


Chosing $\ell=s^{2}$ yields $D \Delta \approx e^{-s^{3 / 2}+s \log d}<1$ for appropriate choice of $s \approx \log ^{2} d$.

