## The complexity of ground states

Zeph Landau

# Ground State

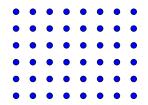
condensed matter

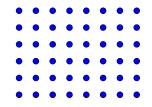
structure gap algorithm

many body Area Law

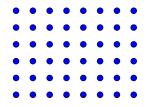
AGSP entanglement viable set

Local Hamiltonian

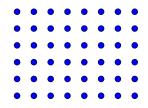




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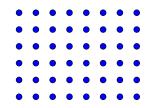


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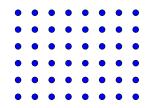
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The same property that leads to the power of quantum computation is the major barrier for understanding many-body physics:

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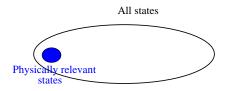
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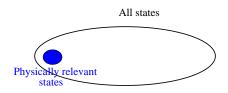
So even describing a state requires exponential amount of information.

### A Basic Question



Can we develop a better understanding of a class of relevant states?

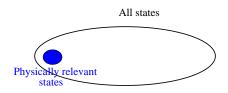
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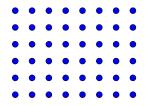
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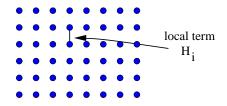
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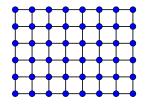
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- acts "locally": non-trivial on only a few particles.

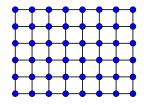


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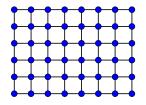
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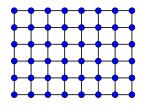
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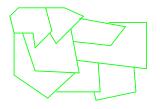
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- Focus on unique ground state and constant gap.

Ground states model the state of the system at low temperatures.

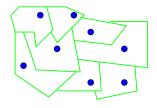
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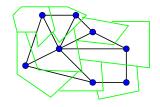
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Each region can be one of three colors,

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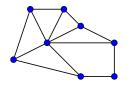
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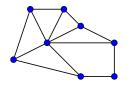
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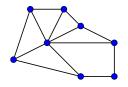


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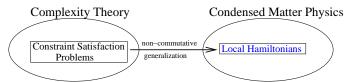
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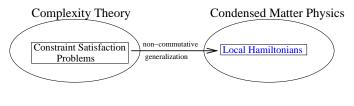


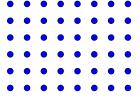
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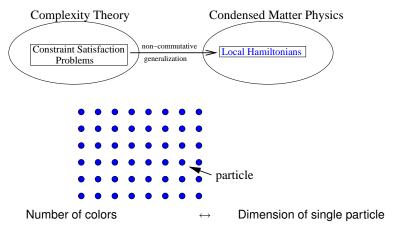
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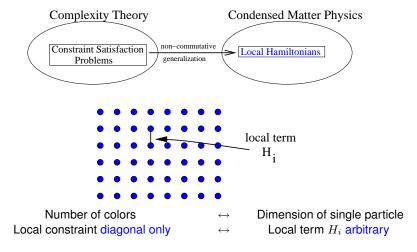
Solving, classifying, and understanding the structure of the solutions of CSP's at the heart of complexity theory.

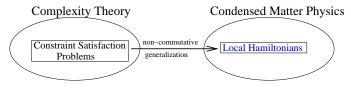


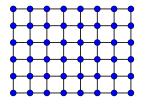












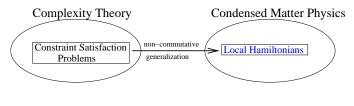
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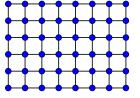
Local constraint diagonal only
Assignment that violates fewest constraints
Least number of constraints violated

Dimension of single particle Local term  $H_i$  arbitrary Ground state: lowest eigenvalue Lowest eigenvalue

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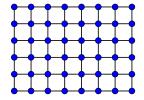
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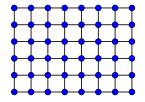
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CSP constraints correspond to  $H_i$  that are diagonal in the standard basis. In particular they all commute.

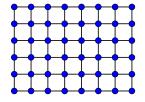
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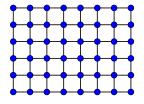




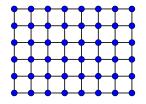
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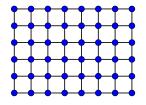
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- For higher dimensions: ?

Understanding ground states of local Hamiltonians: A journey

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- 1D remarkably successful in practice.
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However not a great understanding of what is going on.

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### ['07, Aharonov, Gottesman, Irani, Kempe]

• Solutions to 1D systems are also hard.

### Area Law formulation

Folklore concept motivated by the Holographic Principle in Cosmology:

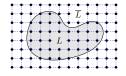
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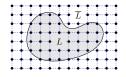
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['01, Vidal, Latorre, Rico, Kitaev] Area Law formalized in terms of entanglement entropy.

 Effect on DMRG: speedup, simplification, better understanding of the heuristics used.



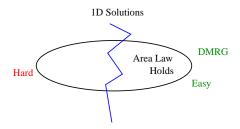
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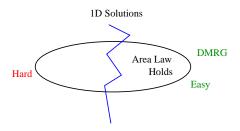
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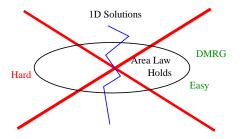


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['08, Cirac, Schuch, Verstraete] Example of finding a solution that satisfies the area law that is NP-hard.

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A special case: frustration-free commuting case.

- Can assume  $H_i$  are projections.
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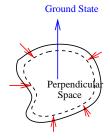
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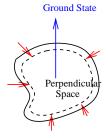
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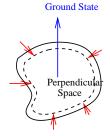
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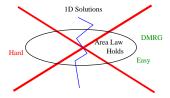
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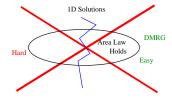


### Properties:

- It "approximately" projects onto one vector you want (ground state).
- It isn't too complex.

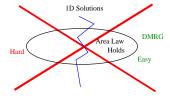






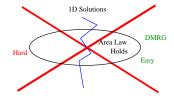
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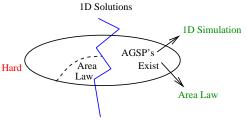
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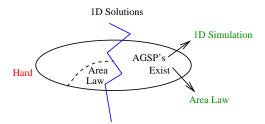


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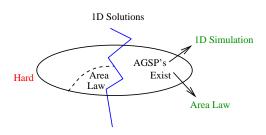
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### Current view of 1D local Hamiltonians



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#### From AGSP's:

- exponential improvement on the constants for the 1D Area Law
- algorithm for 1D,
- gives insight as to what is going on,
- tools for attacking the 2D questions.

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Both steps use AGSPs- the first is much more delicate.

# Measure of Complexity: Entanglement rank

A state on  $\mathcal{H}_1 \otimes \mathcal{H}_2$  of the form  $\sum_{i=1}^{C} a_i \otimes b_i$  will be said to have entanglement rank C.



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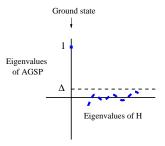
### Entanglement rank behavior

- Multiplicative for operators applied to states or product of operators.
- Additive for sums of states or operators.

We are looking for an operator K with 2 properties:

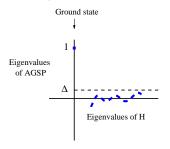
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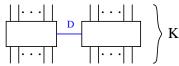


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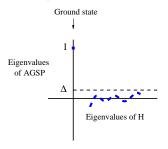


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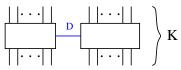


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Critical threshold  $D\Delta < 1$ .

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AGSP definition uses notions: truncation away from cut, Chebyshev polynomials. Requires careful analysis to bound complexity. (See arxiv, find me, or future workshop).

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solving a convex program

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1

#### solving a convex program

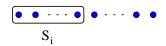
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### Exponential size space is too costly. What we'll need:

- A restriction of the convex program to a polynomial size subspace,
- A succinct description of the elements of that subspace that allows us to perform linear algebra efficiently.

A sequence of spaces  $S_i$  termed viable sets:

- all polynomial size
- all with succinct descriptions that allow efficient linear algebra,
- each containing a good approximation of the "left" side of the ground state.



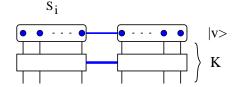
$$|\Gamma\rangle pprox \sum_{j} |a_{j}\rangle \, |b_{j}\rangle \,$$
 with each  $|a_{j}\rangle \in S_{i}.$ 

### Key components:

 Splitting: 1D structure allows reduction of convex program to left side by iterating over a net of boundary conditions. Structural result allows this net to be fixed polynomial size.

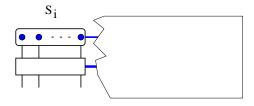
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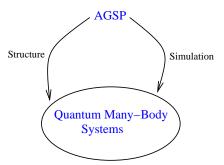


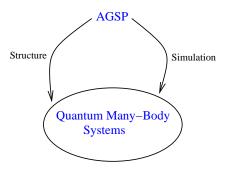
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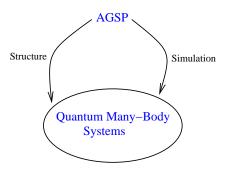


Arxiv, find me, later workshop.

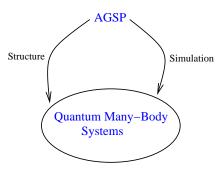




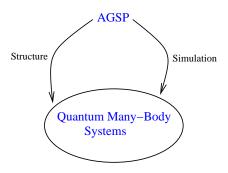
A 2D area law?



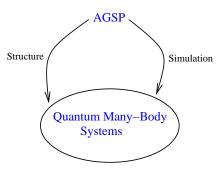
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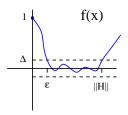


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"The future ain't what it used to be." - Yogi Berra

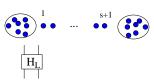
### AGSP construction: norm reduction

Looking for low entanglement operators that look like:



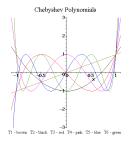
Smaller ||H|| would be better but we don't want to lost the local structure around the cut.

**Solution**: Replace  $H = \sum_i H_i$  with  $H' = H_L + H_1 + H_2 + \cdots + H_s + H_R$ .

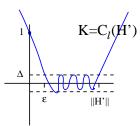


### AGSP construction: Chebyshev polynomials

Chebyshev polynomials: small in an interval:

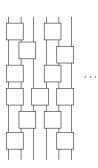


The desired AGSP is a dilation and translation of the Chebyshev polynomial:



$$(H')^{\ell} = \sum (\text{ product of } H_j).$$

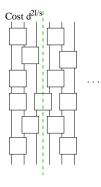
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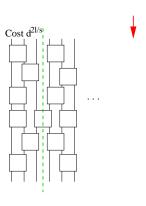
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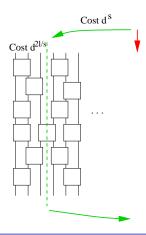
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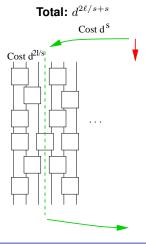
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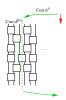
Need to group terms in a nice way but it all works out with total entanglement increase of the same order as the single term.

# Putting things together: Area Law for $H^\prime$

Chebyshev  $C_{\ell}(H')$  has  $\Delta \approx e^{-O(\ell/\sqrt{s})}$ :



Entanglement analysis yields  $D \approx O(d^{\ell/s+s})$ .



Chosing  $\ell=s^2$  yields  $D\Delta\approx e^{-s^{3/2}+s\log d}<1$  for appropriate choice of  $s\approx\log^2 d$ .